

# Relatively Free Algebras with the Identity $x^3 = 0$

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## Abstract

A basis for a relatively free associative algebra with the identity  $x^3 = 0$  over a field of an arbitrary characteristic is found. As an application, a minimal generating system for the  $3 \times 3$  matrix invariant algebra is determined.

## 1 Introduction

Let  $K$  be an infinite field of an arbitrary characteristic  $p$  ( $p = 0, 2, 3, \dots$ ). Let  $K\langle x_1, \dots, x_d \rangle^\#$  ( $K\langle x_1, \dots, x_d \rangle$ , respectively) be the free associative  $K$ -algebra without unity (with unity, respectively) which is freely generated by  $x_1, \dots, x_d$ . Let  $\text{id}\{f_1, \dots, f_s\}$  be the ideal generated by  $f_1, \dots, f_s$ . Denote by  $N_{n,d} = K\langle x_1, \dots, x_d \rangle^\# / \text{id}\{x^n | x \in K\langle x_1, \dots, x_d \rangle^\#\}$  a relatively free finitely generated non-unitary  $K$ -algebra with the identity  $x^n = 0$ , where  $d \geq 1$ . Let  $\mathcal{N} = \{0, 1, 2, \dots\}$ . The algebra  $N_{n,d}$  possesses natural  $\mathcal{N}$ - and  $\mathcal{N}^d$ -gradings by degrees and multidegrees respectively.

The *nilpotency degree* of a non-unitary algebra  $A$  is the least  $C > 0$  for which  $a_1 \cdots a_C = 0$  for all  $a_1, \dots, a_C \in A$ . Denote by  $C(n, d, K)$  the nilpotency degree of  $N_{n,d}$ . In the case of characteristic zero  $n(n+1)/2 \leq C(n, d, K) \leq n^2$  (see [8], [11]), and there is a conjecture that  $C(n, d, K) = n(n+1)/2$ . This conjecture has been proven for  $n \leq 4$  (see [13]). If  $p = 0$  or  $p > n$ , then  $C(n, d, K) < 2^n$  by [6]. For a positive characteristic some upper bounds on  $C(n, d, K)$  are given in [7]:  $C(n, d, K) < (1/6)n^6 d^n$  and  $C(n, d, K) < 1/(m-1)! n^{n^3} d^m$ , where  $m = \lfloor n/2 \rfloor$ .

In [9]  $C(3, d, K)$  was established for an arbitrary  $d, p$ , except for the case of  $p = 3$ ,  $d$  is odd, where the deviation in the estimation of  $C(3, d, K)$  is equal to 1. In this article, a basis for  $N_{3,d}$  is found (see Proposition 2 and Theorems 2, 3), and, in particular,  $C(3, d, K)$  is established for any  $d, p$ . Namely, when  $d > 1$ , we have:

$$\begin{aligned} &\text{If } p = 0 \text{ or } p > 3, \text{ then } C(3, d, K) = 6. \\ &\text{If } p = 2, \text{ then } C(3, d, K) = \begin{cases} d + 3 & , \quad d \geq 3 \\ 6 & , \quad d = 2. \end{cases} \\ &\text{If } p = 3, \text{ then } C(3, d, K) = 3d + 1. \end{aligned}$$

As an application, a minimal homogeneous generating system of the  $3 \times 3$  matrix invariant algebra is determined (see Theorem 4).

For  $p = 2, 3$ , a basis for the multilinear homogeneous component of  $N_{3,d}$  for 'small'  $d$  was found by means of a computer programme. Then, the case of an arbitrary  $d$  was reduced to the case of 'small'  $d$  using the composition method. All programmes were written by means of Borland C++ Builder (version 6.0) and are available upon request from the author. The notion of the composition method was taken from [2].

## 2 Preliminaries

Further, we assume that  $n = 3$ , unless it is stated otherwise. Let  $\mathcal{Z}$  be the ring of integers, and let  $\mathcal{Q}$  be the field of rational fractions. Denote by  $F_d$  the free semigroup, generated by letters  $\{x_1, x_2, \dots, x_d\}$ . By  $F_d^\#$  we mean  $F_d$  without unity. For short, we will write  $K\langle F_d \rangle$  instead of  $K\langle x_1, \dots, x_d \rangle$ . The degree of a  $\mathcal{N}^d$ -homogeneous element  $u \in F_d$  we denote by  $\deg(u)$ , its multidegree we denote by  $\text{mdeg}(u)$ , and the degree of  $u$  in letter  $x_j$  we denote by  $\deg_{x_j}(u)$ . Elements of  $F_d$  are called words. By words from  $N_{3,d}$  we mean images of words from  $F_d$  in  $N_{3,d}$  under the natural homomorphism. We assume that all words are non-empty, that is they are not equal to unity of  $F_d$ , unless it is stated otherwise. Notation  $w = x_{i_1} \cdots \tilde{x}_{i_s} \cdots x_{i_t}$  stands for the word  $w$ , which can be get from the word  $x_{i_1} \cdots x_{i_t}$  by eliminating the letter  $x_{i_s}$ . For a set of words  $M$  and a word  $v$  denote by  $vM$  the set  $\{vu \mid u \in M\}$ . If the set  $M$  is empty, then we assume  $vM = \emptyset$ . By  $x_i \cdots x_j$  ( $i, j \in \mathcal{Z}$ ) we mean the word  $x_i x_{i+1} x_{i+2} \cdots x_j$  if  $1 \leq i \leq j$ , and the empty word otherwise.

For some  $\mathcal{N}^d$ -graded algebra  $A$  and multidegree  $\Delta = (\delta_1, \dots, \delta_d)$  denote by  $A(\Delta)$  or  $A(\delta_1, \dots, \delta_d)$  the homogeneous component of  $A$  of multidegree  $\Delta$ . For short, multidegree  $(3, \dots, 3, 2, \dots, 2, 1, \dots, 1)$  will be denoted by  $3^r 2^s 1^t$  for the appropriate

$r, s, t$ . For  $\Delta = (\delta_1, \dots, \delta_t)$  let  $|\Delta| = \sum_i \delta_i$ . By  $\text{lin}\{v_1, \dots, v_t\}$  we mean the linear span of the elements  $v_1, \dots, v_t$  of some vector space over  $K$ . We denote some elements of  $K\langle F_d \rangle^\#$  by underlined Latin letters.

Endow the set of words of  $F_d$  with the partial lexicographical order. We put  $x_{i_1}x_{i_2}\cdots x_{i_k} < x_{j_1}x_{j_2}\cdots x_{j_t}$  if we have  $i_1 = j_1, \dots, i_{s-1} = j_{s-1}, i_s < j_s$  for some  $s \geq 1$ . Note that if  $v \in F_d^\#$ , then words  $u$  and  $uv$  are incomparable.

By an *identity* we mean an element of  $K\langle F_d \rangle$ . All identities are assumed to be  $N^d$ -homogeneous, unless the contrary is stated. The multidegree of an identity  $t$  we denote by  $\text{mdeg}(t)$ . An identity  $t$  is said to be an identity of  $N_{3,d}$ , if the image of  $t$  in  $N_{3,d}$  under the natural homomorphism is equal to zero. The zero polynomial is called the *trivial* identity.

For identity  $f = \sum_i \alpha_i u_i$ ,  $\alpha_i \in K$ ,  $u_i \in F_d^\#$ ,  $\bar{f}$  stands for the highest term of  $f$ , i.e.,  $\bar{f}$  is the maximal word from the set  $\{u_i\}$ . It is easy to see that, due to homogeneity, the highest term is unique. For a set of identities  $M$ , denote by  $\overline{M}$  the set of the highest terms of the elements from  $M$ .

An identity is called *reduced* if the coefficient of its highest term is equal to 1.

We say that the identity is a *consequence* of a set of identities, provided it belongs to the linear span of these identities. As an example, we point out that  $x_1^3$  is not a consequence of  $x_2^3$ .

An element  $\sum_i \alpha_i u_i \in K\langle F_d \rangle$ , where  $\alpha_i \in K$ ,  $u_i \in F_d^\#$ , is called an *element generated by words*  $v_1, \dots, v_t$ , if all  $u_i$  are products of some elements from  $\{v_1, \dots, v_t\}$ .

For identities  $t_1, t_2$  and for a set of identities  $M$ , notation  $t_1 = t_2 + \{M\}$  means that  $t_1 \in t_2 + \text{lin } M$ .

Consider an element  $g \in K\langle F_d \rangle$  and an identity  $t = u + (\sum_{i=1}^r \alpha_i u_i)$ , where  $\alpha_i \in K$ ,  $u, u_1, \dots, u_r$  are pairwise different words. Let  $g_1 = g$ . If  $g_k = v_1 u v_2 + \sum_{j=1}^s \beta_j w_j$  for pairwise different words  $v_1 u v_2, w_1, \dots, w_s$  and  $\beta_j \in K$ , then  $g_{k+1} = -\sum_{i=1}^r \alpha_i v_1 u_i v_2 + \sum_{j=1}^s \beta_j w_j$ . Note that  $g_{k+1}$  is not uniquely determined by  $g_k$  and  $t$ . If there is  $k$  such that  $g_k = \sum_{j=1}^s \beta_j w_j$  for some words  $w_1, \dots, w_s$  which do not contain subword  $u$ , then we say that the chain  $g_1, \dots, g_k$  is finite and  $g_k$  is its result. If every chain is finite and they all have one and the same result  $g'$ , then we call the identity  $g'$  the *result of application* of the identity  $t$  with the *marked word*  $u$  to the identity  $g$ . Otherwise we say that the result of application of  $t$  with the marked word  $u$  to  $g$  is indefinite.

The result of substitution  $v_1 \rightarrow u_1, \dots, v_k \rightarrow u_k$  in  $f \in K\langle F_d \rangle$ , where  $f$  is an element generated by words  $v_1, \dots, v_k, v_{k+1}, \dots, v_t$ , denote by  $f|_{v_1 \rightarrow u_1, \dots, v_k \rightarrow u_k}$ . By *substitutional mapping* we mean such homomorphism of  $K$ -algebras  $\phi : K\langle F_k \rangle \rightarrow K\langle F_l \rangle$  that  $\phi(x_i) \in F_l^\#$ ,  $i = \overline{1, k}$ . A substitutional mapping is called monotonous, if

$\phi(x_i) > \phi(x_j)$  for  $x_i > x_j$ . The set of monotonous substitutional mappings denote by  $\mathcal{M}_{k,l}$ . Note that for  $\phi \in \mathcal{M}_{i,j}$ ,  $\psi \in \mathcal{M}_{j,k}$  the composition  $\psi \circ \phi$  belongs to  $\mathcal{M}_{i,k}$ . Denote

$$T_1(a) = a^3,$$

$$T_2(a, b) = a^2b + aba + ba^2,$$

$$T_3(a, b, c) = abc + acb + bac + bca + cab + cba.$$

Partial and complete linearization of the identity  $f_1a^3f_2$  of  $N_{3,d}$ , where  $a \in K\langle F_d \rangle^\#$ ,  $f_1, f_2 \in F_d$ , gives that all identities from  $\mathcal{S} = \{f_1T_1(a)f_2, f_1T_2(a, b)f_2, f_1T_3(a, b, c)f_2 \mid a, b, c \in F_d^\#, f_1, f_2 \in F_d\}$  are identities of  $N_{3,d}$ . For multidegree  $\Delta$  let  $\mathcal{S}_\Delta$  be the subset of  $\mathcal{S}$  which consists of all identities of multidegree  $\Delta$ . Clearly, each identity of  $N_{3,d}(\Delta)$  is a consequence of the set of identities  $\mathcal{S}_\Delta$ . The set of identities  $\mathcal{S}_\Delta$  can be treated like the system of homogeneous linear equations in formal variables  $\{w \mid w \in F_d^\#, \text{mdeg}(w) = \Delta\}$ . Then, free variables of the system  $\mathcal{S}_\Delta$  form a basis for  $N_{3,d}(\Delta)$ . We call two systems of linear equations (two sets of identities, respectively) equivalent, if the first one is a consequence of the second and vice versa.

A word  $w \in S$  is called *canonical* with respect to  $x_i$ , if it has one of the following forms:  $w_1, w_1x_iw_2, w_1x_i^2w_2, w_1x_i^2ux_iw_2$ , where words  $w_1, w_2, u$  do not contain  $x_i$ , words  $w_1, w_2$  can be empty. If a word is canonical with respect to all letters, then we call it *canonical*. Number for future references the identity of  $N_{3,d}$

$$xux + (x^2u + ux^2), \quad u \in F_d^\#. \quad (1)$$

### 3 Auxiliary results

We will use the following facts from [9]:

**Lemma 1** 1. Applying identities (1),  $x_iux_i^2 = -x_i^2ux_i$  of  $N_{3,d}$ , any non-zero word  $w \in N_{3,d}$  can be represented as a sum of canonical words which belong to the same homogeneous component as  $w$ . In particular, if  $\deg_{x_i}(w) > 3$ ,  $w \in F_d$ , then  $w = 0$  in  $N_{3,d}$ .

2. The inequality  $x_1^2x_2^2x_1 \neq 0$  holds in  $N_{3,d}$ .

3. If  $p = 0$  or  $p > 3$ , then  $C(3, d, K) = 6$  ( $d \geq 2$ ).

If  $p = 2$ , then  $C(3, d, K) = d + 3$ , where  $d \geq 3$ , and  $C(3, 2, K) = 6$ .

4. If  $p \neq 3$ , then  $x^2ay^2 = 0$  is an identity of  $N_{3,d}$ , where  $a \in F_d^\#$ .

5. If  $p = 3$ , then  $x^2y^2xay = x^2y^2xya$  is an identity of  $N_{3,d}$ , where  $a \in F_d^\#$ .

6. If  $p = 2$ , then  $x_1^2 x_2 \cdots x_d x_1 \neq 0$  holds in  $N_{3,d}$ .  
7. If  $p \neq 3$ , then  $I_1(x, a, b, c) = x^2 abc + x^2 acb$ ,  $I_2(x, a, b, c) = abc x^2 + bac x^2$ ,  $I_3(x, a, b, c) = ax^2 bc + cax^2 b$  are identities of  $N_{3,d}$ , where  $x, a, b, c$  are words.

**Proof.** 1. See [9], Statement 1.

2. See [9], Statement 3.

3. See [9], Propositions 1, 2.

4. See [9], equality (5).

5. See [9], proof of Statement 7.

6. See [9], Statement 4.

7. Let  $x, a, b, c$  be words,  $p \neq 3$ . Partial linearization of the identity from item 4 with respect to  $x$  ( $y$ , respectively) gives that  $I_1(x, a, b, c)$ ,  $I_2(x, a, b, c)$  are identities of  $N_{3,d}$ . Apply identity (1), where  $x = x_1$ , to the identity  $T_3(x_1 a, b x_1, c) = 0$  of  $N_{3,d}$ , and get that  $-T_3(x_1^2 a, b, c) - T_3(a, b x_1^2, c) + 3(b x_1^2 a c + c b x_1^2 a) = 0$  in  $N_{3,d}$ . Hence  $I_3(x, a, b, c)$  is the identity of  $N_{3,d}$ .  $\triangle$

**Remark 1** 1. Consider a set  $M = \{m_i\}_{i=\overline{1,s}} \subset K\langle F_d \rangle^\#$ . Let  $u \in F_d^\#$  be a word which is a summand of one and only one element  $m_1$  of the set  $M$ . Let  $\sum_{i=1}^s \alpha_i m_i = 0$  in  $K\langle F_d \rangle$ , where  $\alpha_i \in K$ . Then  $\alpha_1 = 0$ .

2. Let  $\Delta$  be a multidegree. Let  $V = \{v_1, \dots, v_s\}$  be a set of words of multidegree  $\Delta$ . Assume that for each word  $w$ , of multidegree  $\Delta$ , which do not lie in  $V$  there is an identity  $w - f_w$  of  $N_{3,d}$ , where  $f_w \in \text{lin } V$ . Then every identity  $\sum_i \alpha_i v_i$  ( $\alpha_i \in K$ ) of  $N_{3,d}(\Delta)$  is a consequence of the identities which are results of application of identities  $\{w - f_w\}$  to the identities of  $\mathcal{S}_\Delta$ . (Note that results of these applications are defined.)

**Lemma 2** Let  $d \geq 1$ . All identities, of  $N_{3,d}(21^{d-1})$ , generated by  $x_1^2, x_2, \dots, x_d$  are consequences of the following identities of  $N_{3,d}(21^{d-1})$ :

- (a)  $f_1 T_3(a, b, c) f_2$ , where some word from  $a, b, c, f_1, f_2$  contains the subword  $x_1^2$ ,  
(b)  $3f_1 I_i(x_1, a, b, c) f_2$ ,  $i = \overline{1, 3}$ .

Here  $a, b, c \in F_d^\#$ ,  $f_1, f_2 \in F_d$ .

**Proof.** By item 2 of Remark 1, any identity, of  $N_{3,d}(21^{d-1})$ , generated by  $x_1^2, x_2, \dots, x_d$  is a consequence of identities which are results of application of identity (1) (where  $x = x_1$ ) to the identities from  $\mathcal{S}_{21^{d-1}}$ .

If we apply (1), where  $x = x_1$ , to  $T_2(x_1, a) = 0$ , then we get a trivial identity.

The result of application of (1), where  $x = x_1$ , to an identity  $t = f_1 T_3(a_1, a_2, a_3) f_2$ , where  $a_1, a_2, a_3 \in F_d^\#$ ,  $f_1, f_2 \in F_d$ , denote by  $t'$ , and let  $t = \sum_{i=1}^6 u_i$  for some words  $u_1, \dots, u_6$ . If words  $u_1, \dots, u_6$  do not contain subword  $x_1^2$  and  $a_i \neq x_1$ ,  $i = \overline{1, 3}$ , then  $t'$  is a consequence of identities (a). Let  $a, b, c \in F_d^\#$ .

If  $t = T_3(x_1, x_1, a)$ , then  $t' = 0$ .

If  $t = T_3(x_1 a, x_1, b)$ , then  $t' = -T_3(x_1^2, a, b)$ .

If  $t = T_3(ax_1 b, x_1, c)$ , then  $t' = -T_3(ab, x_1^2, c) - 3I_3(x_1, a, b, c)$ .

If  $t = T_3(x_1 a, bx_1, c)$ , then  $t' = -T_3(x_1^2 a, b, c) - T_3(a, bx_1^2, c) + 3I_3(x_1, b, a, c)$ .

If  $t = x_1 a T_3(x_1, b, c)$ , then  $t' = -a T_3(x_1^2, b, c) - 3I_1(x_1, a, b, c)$ .

If  $t = x_1 T_3(x_1, a, b)$ , then  $t' = -T_3(x_1^2, a, b)$ .

If  $t = x_1 T_3(x_1 a, b, c)$ , then  $t' = -x_1^2 T_3(a, b, c) - T_3(x_1^2 a, b, c) + 3I_1(x_1, a, b, c)$ .

Due to the fact that, if we read identities (a), (b) from right to left, they do not change, the claim follows from the regarded cases.  $\triangle$

Let  $r = \overline{1, d}$ . It is easy to see that for every  $i = \overline{1, r}$  the result of application of the identity (1), where  $x = x_i$ , to every identity of multidegree  $2^r 1^{d-r}$  is definite. For  $\sigma \in S_r$  let  $\psi_\sigma : K\langle F_d \rangle(2^r 1^{d-r}) \rightarrow K\langle F_d \rangle(2^r 1^{d-r})$  be such mapping that  $\psi_\sigma(t)$  is the result of the following procedure. Let  $t_1$  be the result of application of the identity (1), where  $x = x_{\sigma(1)}$ , to  $t$ . For  $i = \overline{2, r}$  let  $t_i$  be the result of application of the identity (1), where  $x = x_{\sigma(i)}$ , to  $t_{i-1}$ . We define  $\psi_\sigma(t) = t_r$ .

For any identity  $t = \sum_{i=1}^s \alpha_i u_i$ ,  $\alpha_i \in K$ ,  $u_i \in F_d^\#$ , of multidegree  $2^r 1^{d-r}$ , fix some permutations  $\sigma_{t,1}, \dots, \sigma_{t,s} \in S_r$ . Consider the mapping  $\psi : K\langle F_d \rangle(2^r 1^{d-r}) \rightarrow K\langle F_d \rangle(2^r 1^{d-r})$  such that for  $t = \sum_i \alpha_i u_i$ ,  $\alpha_i \in K$ ,  $u_i \in F_d^\#$ , we have  $\psi(t) = \sum_i \alpha_i \psi_{\sigma_{t,i}}(u_i)$ . Denote by  $\Psi_r$  the set of all such mappings  $\psi$ .

**Lemma 3** *Let  $d, r \geq 1$ ,  $\phi \in \Psi_r$ . All identities, of  $N_{3,d}(2^r 1^{d-r})$ , generated by  $x_1^2, \dots, x_r^2, x_{r+1}, \dots, x_d$  are consequences of the following identities of  $N_{3,d}(2^r 1^{d-r})$ :*

(a)  $f_1 T_3(a, b, c) f_2$ , where for each  $k = \overline{1, r}$  some word from  $a, b, c, f_1, f_2$  contains the subword  $x_k^2$ ,

(b)  $3\phi(f_1 I_i(x_k, a, b, c) f_2)$ ,  $i = \overline{1, 3}$ ,  $k = \overline{1, r}$ ,

(c) the identity  $3f_1 x_i^2 a x_j^2 f_2$  ( $i, j = \overline{1, r}$ ,  $i \neq j$ ) which is generated by  $x_1^2, \dots, x_r^2, x_{r+1}, \dots, x_d$ .

Here  $a, b, c \in F_d^\#$ ,  $f_1, f_2 \in F_d$ .

**Proof.** For  $\psi \in \Psi_r$  denote the sets of identities of multidegree  $2^r 1^{d-r}$ :  $A_1 = \{\psi_\sigma(w) - \psi_\tau(w) \mid w \in F_d, \sigma, \tau \in S_r\}$ ,  $A_2^\psi = \{\psi(t) \mid t = f_1 T_2(a, b) f_2, a, b \in F_d^\#, f_1, f_2 \in F_d\}$  and  $A_3^\psi = \{\psi(t) \mid t = f_1 T_3(a, b, c) f_2, a, b, c \in F_d^\#, f_1, f_2 \in F_d\}$ .

For an arbitrary  $\psi \in \Psi_r$  identities, of  $N_{3,d}(2^r 1^{d-r})$ , generated by  $x_1^2, \dots, x_r^2, x_{r+1}, \dots, x_d$  are consequences of  $A_1, A_2^\psi, A_3^\psi$  (see item 2 of Remark 1). The following items conclude the proof.

1. *Identities  $A_1$  are consequences of identities (c). In particular,  $\psi(t) = \pi(t) + \{(a), (b), (c)\}$  and  $\psi(t + f) = \psi(t) + \psi(f)$  for any  $\psi, \pi \in \Psi_r$ ,  $t, f \in K\langle F_d \rangle(2^r 1^{d-r})$ .*

Proof. Note that the identity  $\psi(3f_1x_i^2ax_j^2f_2) \in K\langle F_d \rangle(2^r1^{d-r})$ , where  $a \in F_d^\#$ , follows from (c).

If  $r = 1$  then  $A_1 = \{0\}$ . Let  $r = 2$ . Denote  $\psi_1(w) = \psi_\sigma(w) - \psi_\tau(w)$ , where  $\sigma = 1 \in S_2$ ,  $\tau = (1, 2) \in S_2$ . If  $w = x_1ax_1bx_2cx_2$  or  $w = x_1ax_2bx_2cx_1$ ,  $a, b, c \in F_d$ , then the identity  $\psi_1(w)$  is trivial. Consider  $w = x_1ax_2bx_1cx_2$ ,  $a, b, c \in F_d$ . Then  $\psi_1(w) = (x_1ax_2bx_1)cx_2 - x_1a(x_2bx_1cx_2)$  in  $N_{3,d}$ , where the order of application of identity (1) is determined by parentheses.

If  $ab, bc \neq 1$ , then  $\psi_1(w) = (-x_1^2ax_2bcx_2 - ax_2bx_1^2cx_2) - (-x_1ax_2^2bx_1c - x_1abx_1cx_2^2) = (x_1^2ax_2^2bc + x_1^2abcx_2^2 + ax_2^2bx_1^2c + abx_1^2cx_2^2) - (x_1^2ax_2^2bc + ax_2^2bx_1^2c + x_1^2abcx_2^2 + abx_1^2cx_2^2) = 0$ .

If  $a = b = c = 1$ , then  $\psi_1(w) = 0$ .

If  $a = b = 1$ ,  $c \neq 1$ , then  $\psi_1(w) = 3x_1^2cx_2^2$ .

If  $b = c = 1$ ,  $a \neq 1$ , then  $\psi_1(w) = -3x_1^2ax_2^2$ .

Therefore, if  $r = 2$  then the required is proved.

The case of  $r > 2$  follows from the case of  $r = 2$  and the fact that any permutation is a composition of elementary transpositions.

2. Identities  $A_3^\psi$  are consequences of identities (a), (b), (c).

Proof. It follows from Lemma 2 and item 1.

3. Identities  $A_2^\psi$  are consequences of identities (a), (b), (c).

Proof. We will use item 1 without reference. Prove by induction on  $k$  that for every identity  $t = f_1T_2(v, a)f_2$  of multidegree  $2^r1^{d-r}$  we have  $\psi(t) = 0 + \{(a), (b), (c)\}$ , i.e.,  $\psi(f_1 \cdot vav \cdot f_2) = -\psi(f_1 \cdot v^2a \cdot f_2) - \psi(f_1 \cdot av^2 \cdot f_2) + \{(a), (b), (c)\}$ , where  $v, a \in F_d^\#$ ,  $f_1, f_2 \in F_d$ , and  $\deg(v) = k$ .

Induction base is trivial.

Induction step. Without loss of generality we can assume that  $f_1, f_2$  are empty words. Consider a word  $x_iu$  of degree  $k$ , where  $i = \overline{1, r}$ . We have  $\psi(T_2(x_iu, a)) = \psi(x_iux_iua) + \psi(ax_iux_iu) + \psi(x_iuax_iu) + \{(a), (b), (c)\}$ . Induction hypothesis imply that  $\psi(x_iux_iua) = \psi(u^2x_i^2a) + \{(a), (b), (c)\}$ ,  $\psi(ax_iux_iu) = \psi(au^2x_i^2) + \{(a), (b), (c)\}$ ,  $\psi(x_iuax_iu) = \psi(x_i^2u^2a) + \psi(x_i^2au^2) + \psi(u^2ax_i^2) + \psi(ax_i^2u^2) + \{(a), (b), (c)\}$ . Thus,  $\psi(T_2(x_iu, a)) = \psi(T_3(x_i^2, u^2, a)) + \{(a), (b), (c)\} = 0 + \{(a), (b), (c)\}$  by item 2.  $\triangle$

**Lemma 4** Let  $p = 3$ . All identities of  $N_{3,d}(1^d)$  are consequences of identities  $f_1T_3(a_1, a_2, a_3)f_2$  of multidegree  $1^d$ , where  $f_1, f_2 \in F_d$ ,  $a_1, a_2, a_3 \in F_d^\#$ ,  $\deg(a_1) \leq 3$ ,  $\deg(a_2) = \deg(a_3) = 1$ .

**Proof.** If  $d \leq 4$ , then the statement is obvious.

Let  $d \geq 5$ . We prove by induction on  $d$ .

Induction base. In the case  $d = 5, 6$  the statement was proven by means of a computer programme.

Induction step. Let  $d \geq 7$ . Consider an identity  $t = a_1 T_3(a_2, a_3, a_4) a_5$ ,  $a_1, a_5 \in F_d$ ,  $a_2, a_3, a_4 \in F_d^\#$ . There is such  $k = \overline{1, 5}$  that  $\deg(a_k) \geq 2$ . Then  $a_k = x_i x_j \cdot w$  for some  $w \in F_d$ . Substituting  $z$  for subword  $x_i x_j$ , where  $z$  is a new letter, and using induction hypothesis, we get that  $t \in \text{lin}\{f_1 T_3(b_1, b_2, b_3) f_2 \mid f_1, f_2 \in F_d, b_1, b_2, b_3 \in F_d^\#, \deg(b_1 b_2 b_3) \leq 6\}$ . The statement of the Lemma in the case  $d = 6$  concludes the proof.  $\triangle$

A multilinear word  $w = x_{\sigma(1)} \cdots x_{\sigma(d)}$ ,  $\sigma \in S_d$ , is called *even*, if the permutation  $\sigma$  is even, and *odd* otherwise.

**Lemma 5** 1. Let  $p = 2, 3$ . Consider the homomorphism  $\phi : K\langle F_d \rangle(1^d) \rightarrow K$ , defined by  $\phi(w) = 1$ , where  $w \in F_d$ . Then  $\phi$  maps identities of  $N_{3,d}$  in zero.

2. Let  $p = 2$ . For  $i, j = \overline{1, d}$ ,  $i \neq j$ , consider the homomorphism  $\phi_{ij} : K\langle F_d \rangle(1^d) \rightarrow K$ , defined by the following way: if  $w = u x_i x_j v$ , where  $u, v \in F_d$ , then  $\phi_{ij}(w) = 1$ , else  $\phi_{ij}(w) = 0$ . Then  $\phi_{ij}$  maps identities of  $N_{3,d}$  in zero.

3. Let  $p = 3$ . Consider the homomorphism  $\phi_+ : K\langle F_d \rangle(1^d) \rightarrow K$ , defined by the following way: for  $w \in F_d$  we have  $\phi_+(w) = 1$ , if  $w$  is even, else  $\phi_+(w) = 0$ . Then  $\phi_+$  maps identities of  $N_{3,d}$  in zero.

4. Let  $p = 3$ . For  $k = \overline{1, d}$  consider the homomorphism  $\phi_k : K\langle F_d \rangle(\delta_1, \dots, \delta_d) \rightarrow K\langle F_d \rangle(\delta_1, \dots, \delta_{k-1}, \delta_{k+1}, \dots, \delta_d)$ , where  $\delta_k = 1, 2$ , defined by  $\phi_k(x_{i_1} \cdots x_{i_t}) = y_{i_1} \cdots y_{i_t}$ , where  $y_{i_j} = x_{i_j}$ , if  $i_j \neq k$ , and  $y_{i_j} = 1$ , if  $i_j = k$  ( $j = \overline{1, t}$ ,  $t = \delta_1 + \cdots + \delta_d$ ). Then  $\phi_k$  maps identities of  $N_{3,d}$  in zero.

5. Let  $p = 3$ . For  $k = \overline{1, d}$  consider the homomorphism  $\pi_k : K\langle F_d \rangle(\delta_1, \dots, \delta_{k-1}, 3, \delta_{k+1}, \dots, \delta_d) \rightarrow K\langle F_d \rangle(\delta_1, \dots, \delta_{k-1}, 1, \delta_{k+1}, \dots, \delta_d)$ , defined by the following way: for  $w = u_1 x_k u_2 x_k u_3 x_k u_4$ ,  $u_i \in F_d$  ( $i = \overline{1, 4}$ ), we put  $\pi_k(w) = u_1(x_k u_2 u_3 + u_2 x_k u_3 + u_2 u_3 x_k) u_4$ . Then  $\pi_k$  maps identities of  $N_{3,d}$  in zero, and  $\pi_i \pi_j(x_i^2 x_j^2 x_i x_j) = x_i x_j - x_j x_i$ .

**Proof.** 2. It is sufficient to proof that for  $t = b_1 T_3(a_1, a_2, a_3) b_2$ , where  $b_1, b_2 \in F_d$ ,  $a_1, a_2, a_3 \in F_d^\#$ , we have  $\phi_{ij}(t) = 0$ . If each word from  $\{b_1 a_{\sigma(1)} a_{\sigma(2)} a_{\sigma(3)} b_2 \mid \sigma \in S_3\}$  do not contain subword  $x_i x_j$ , then  $\phi_{ij}(t) = 0$ .

If there is  $k$  such that  $a_k = u x_i x_j v$ ,  $u, v \in F_d$ , then  $\phi_{ij}(t) = 6 = 0$ .

If  $b_1 = u x_i$ ,  $a_1 = x_j v$ ,  $u, v \in F_d$ , then  $\phi_{ij}(t) = 2 = 0$ .

If  $b_2 = x_j u$ ,  $a_1 = v x_i$ ,  $u, v \in F_d$ , then  $\phi_{ij}(t) = 2 = 0$ .

If  $a_1 = u x_i$ ,  $a_2 = x_j v$ ,  $u, v \in F_d$ , then  $\phi_{ij}(t) = 2 = 0$ .

The statement follows from the regarded cases.

Items 3, 4, 5 were proved in [9]; item 1 is similar to them.  $\triangle$



## 4 The case of $p \neq 3$ , $d \leq 3$

**Proposition 1** *Let  $p \neq 3$ ,  $d = \overline{1, 3}$ ,  $\Delta = (\delta_1, \dots, \delta_d)$  is a multidegree. Then, the set  $B_\Delta$  is a basis for  $N_{3,d}(\Delta)$ , where*

1) *the case of  $|\Delta| \leq 3$ :*

$$B_1 = \{x_1\}, B_{1^2} = \{x_1x_2, x_2x_1\}, B_2 = \{x_1^2\},$$

$$B_{1^3} = \{x_1x_2x_3, x_1x_3x_2, x_2x_1x_3, x_2x_3x_1, x_3x_1x_2\}, B_{21} = \{x_1^2x_2, x_2x_1^2\};$$

2) *the case of  $|\Delta| = 4$ :*

$$B_{21^2} = \{x_1^2x_2x_3, x_1^2x_3x_2, x_2x_1^2x_3, x_2x_3x_1^2, x_3x_2x_1^2\},$$

$$B_{2^2} = \{x_1^2x_2^2, x_2^2x_1^2\}, B_{31} = \{x_1^2x_2x_1\};$$

3) *the case of  $|\Delta| = 5$ :*

$$B_{2^21} = \{x_1^2x_2^2x_3, x_2^2x_1^2x_3, x_3x_1^2x_2^2\},$$

$$B_{31^2} = \{x_1^2x_2x_3x_1, x_1^2x_3x_2x_1\},$$

$$B_{32} = \{x_1^2x_2^2x_1\};$$

4) *otherwise  $B_\Delta = \emptyset$ .*

**Proof.** Cases of  $|\Delta| \leq 3$  and  $\Delta \in \{31, 32\}$  follow from items 1, 2 of Lemma 1.

If  $\Delta \in \{21^2, 2^2, 2^21, 31^2\}$ , then we prove the statement by considering the system of equations  $\mathcal{S}_\Delta$ . Here we use item 1 of Lemma 1; and when  $\Delta \in \{21^2, 2^2, 2^21\}$ , we use Lemma 3 for decreasing the number of considering equations.

Case 4) follows from item 3 of Lemma 1.  $\Delta$

## 5 The case of $p = 0$ or $p > 3$

We shall write  $i$  for  $x_i$ ,  $i = \overline{1, d}$ , so that it does not lead to ambiguity.

**Proposition 2** *Let  $p = 0$  or  $p > 3$ ,  $d \geq 1$ ,  $\Delta = (\delta_1, \dots, \delta_d)$  is a multidegree. Then the set  $B_\Delta$  is a basis for  $N_{3,d}(\Delta)$ , where*

1) *if  $d \leq 3$ , then see Proposition 1,*

$$2) B_{1^4} = \left\{ \begin{array}{l} 1234, 1243, 1324, 1342, 1423, \\ 2134, 2143, 2314, 2341, 2413, \\ 3124, 3412 \end{array} \right\},$$

$$3) B_{1^5} = \left\{ \begin{array}{l} 12345, 12354, 12435, 12453, 12534, \\ 13245, 13254, 13425, 13452, 13524, \\ 14235, 14523, 23145, 23415, 23514 \end{array} \right\},$$

$$4) B_{21^3} = \left\{ \begin{array}{l} 1^2234, 1^2324, 1^2423, \\ 21^234, 21^243, 231^24, 241^23 \end{array} \right\},$$

5) *if  $|\Delta| \geq 6$ , then  $B_\Delta = \emptyset$ .*

**Proof.** The computations below were performed by means of a computer programme. For  $\Delta \in \{1^4, 1^5, 21^4\}$ , we consider the homogeneous system of linear equations  $\mathcal{S}_\Delta$  over the ring, generated in  $\mathcal{Q}$  by the set  $\mathcal{Z} \cup \{1/2, 1/3\}$ . Having expressed higher words in terms of lower words by the Gauss's method, we get that  $\mathcal{S}_\Delta$  is equivalent to the system  $\{1 \cdot u = f_u \mid u \in F_d, \text{mdeg}(u) = \Delta, u \notin B_\Delta\}$ , where  $f_u$  are linear combinations of elements of  $B_\Delta$ . The statement is proven.

The case of  $|\Delta| \geq 6$  follows from Lemma 1.  $\triangle$

## 6 The composition method

Denote by  $M_5$  a basis for the space of identities of  $N_{3,5}(1^5)$  such that  $M_5$  contains only reduced identities and all elements of  $M_5$  have the highest terms pairwise different. The basis of this kind exists, because if  $t_1, t_2$  are reduced identities with  $\bar{t}_1 = \bar{t}_2$  and  $t_1 \neq t_2$ , then  $\bar{t}_1 \neq \overline{t_1 - t_2}$  and  $\text{lin}\{t_1, t_2\} = \text{lin}\{t_1, t_1 - t_2\}$ . For  $d \geq 5$ , let

$$M_d = \{t \in K\langle F_d \rangle(1^d) \mid \text{there are } t' \in M_5, \phi \in \mathcal{M}_{5,d}, a \in F_d \text{ such that } t = a\phi(t')\}.$$

Identities from  $M_d$  are identities of  $N_{3,d}$ .

Let  $\Delta = (\delta_1, \dots, \delta_d)$  be a multidegree. For a set  $J \subset K\langle F_d \rangle(\Delta)$ , denote  $B(J) = \{w \in F_d^\# \mid \text{mdeg}(w) = \Delta, w \notin \overline{J}\}$ . Since every word, of multidegree  $\Delta$ , which do not belong to  $B(J)$ , can be expressed in terms of lower words by applying identities of  $J$ ; therefore, for any  $f \in K\langle F_d \rangle(\Delta)$ , we have

$$f = \sum_i \alpha_i t_i + \sum_j \beta_j w_j, \text{ where } \alpha_i, \beta_j \in K, t_i \in J, w_j \in B(J), \bar{t}_i, w_j \leq \bar{f}. \quad (2)$$

Thus,

$$\frac{K\langle F_d \rangle(\Delta)}{\text{lin } J} \simeq \text{lin } B(J), \text{ and, in particular, } N_{3,d}(1^d) = \text{lin } \Phi(B(M_d)), \quad (3)$$

where  $\Phi : K\langle F_d \rangle^\# \rightarrow N_{3,d}$  is the natural homomorphism. Further, we will write  $B(M_d)$  instead of  $\Phi(B(M_d))$  so that it does not lead to ambiguity.

**Definition.** A set of reduced identities  $M$  is called *complete under composition*, if for any  $t_1, t_2 \in M$ , where  $\bar{t}_1 = \bar{t}_2$ , we have  $t_1 - t_2 = \sum_{i=1}^k \alpha_i g_i$ ,  $\alpha_i \in K$ ,  $g_i \in M$ , and  $\bar{g}_i < \bar{t}_1$ ,  $i = \overline{1, k}$ .

**Lemma 6** For  $d \geq 5$ , all identities of  $N_{3,d}(1^d)$  are consequences of identities of  $M_d$ .

**Proof.** For any identity  $t$  of  $N_{3,d}(1^d)$ , we have  $t = \sum_i \alpha_i t_i$  for some  $\alpha_i \in K$ ,  $t_i = f_i T_3(a_i, b_i, c_i) g_i$ ,  $f_i, g_i \in F_d$ ,  $a_i, b_i, c_i \in F_d^\#$ . For any  $i$ , exists  $\phi_i \in \mathcal{M}_{5,d}$  and an identity  $t'_i$  of  $N_{3,5}(1^5)$ , such that  $t_i = \phi_i(t'_i)$ . The set  $M_5$  is a basis for the space of identities of  $N_{3,5}(1^5)$ , thus  $t'_i = \sum_j \beta_{ij} g_{ij}$ , where  $g_{ij} \in M_5$ ,  $\beta_{ij} \in K$ . Hence  $\phi_i(t'_i) = \sum_j \beta_{ij} \phi_i(g_{ij}) \in \text{lin } M_d$ . Therefore,  $t \in \text{lin } M_d$ .  $\triangle$

**Lemma 7** (*Composition Lemma [2]*) For  $d \geq 5$   $B(M_d)$  is a basis for  $N_{3,d}(1^d)$  if and only if  $M_d$  is complete under composition.

**Proof.** Let  $B(M_d)$  be a basis for  $N_{3,d}(1^d)$ . For  $t_1, t_2 \in M_d$ , where  $\bar{t}_1 = \bar{t}_2$ , let  $g = t_1 - t_2$ . By formula (2) we have  $g = \sum_i \alpha_i f_i + \sum_j \beta_j w_j$ , where  $\bar{f}_i \leq \bar{g} < \bar{t}_1$ ,  $f_i \in M_d$ ,  $w_i \in B(M_d)$ ,  $\alpha_i, \beta_j \in K$ . The identity  $g - \sum_i \alpha_i f_i$  is an identity of  $N_{3,d}$ ,  $B(M_d)$  is a basis for  $N_{3,d}(1^d)$ , hence  $\sum_j \beta_j w_j = 0$  in  $K\langle F_d \rangle$ . Thus  $g = \sum_i \alpha_i f_i$ , and the claim is proven.

Let  $M_d$  be complete under composition. Assume that, on the contrary,  $B(M_d)$  is not a basis. Thus, the set  $B(M_d)$  is linearly dependent in  $N_{3,d}$  (see equality (3)). Hence there is a non-trivial identity  $f = \sum_i \alpha_i u_i$ ,  $\alpha_i \in K$ ,  $u_i \in B(M_d)$ , such that  $f = 0$  in  $N_{3,d}$ . Note that  $\bar{f} \notin \overline{M}_d$ .

Lemma 6 implies  $f = \sum_{j=1}^k \beta_j t_j$ , where  $\beta_j \in K^*$ ,  $t_j \in M_d$ . Without loss of generality, we can assume that for some  $s$ , we have  $\bar{t}_1 = \dots = \bar{t}_s = \bar{f}$ . If  $s = 1$  then we get a contradiction to  $\bar{f} \notin \overline{M}_d$ . Let  $s \geq 2$ . Since  $M_d$  is complete under composition, for  $j = \overline{2, s}$  we have  $t_1 - t_j = \sum_l \gamma_{jl} g_{jl}$ , where  $g_{jl} \in M_d$ ,  $\bar{g}_{jl} < \bar{t}_1$ ,  $\gamma_{jl} \in K$ . Expressing  $t_j$  from these equalities, we get  $f = \lambda t_1 + \sum_q \lambda_q h_q$  for some  $h_q \in M_d$ ,  $\bar{h}_q < \bar{t}_1$ ,  $\lambda, \lambda_q \in K$ . If  $\lambda \neq 0$ , then  $\bar{f} \in \overline{M}_d$ , so we get a contradiction. Thus,  $\lambda = 0$ . Repeating the same argument several times, we get a contradiction to the non-triviality of  $f$ .  $\triangle$

**Lemma 8** Let  $d \geq 5$ . Then for any  $t_1, t_2 \in M_d$ , where  $\bar{t}_1 = \bar{t}_2$ , there are  $s = \overline{5, 10}$ ,  $t'_1, t'_2 \in M_s$ ,  $\phi \in \mathcal{M}_{s,d}$ ,  $a \in F_d$  such that  $t_i = a\phi(t'_i)$ , where  $i = 1, 2$ .

**Proof.** By definition of  $M_d$ , there are  $t''_i \in M_5$ ,  $\psi_i \in \mathcal{M}_{5,d}$ ,  $c_i \in F_d$  such that  $t_i = c_i \psi_i(t''_i)$ , where  $i = 1, 2$ . Denote  $w = \bar{t}_1 = \bar{t}_2$ . Let  $\bar{t}''_1 = x_{j_1} \dots x_{j_5}$ ,  $\bar{t}''_2 = x_{k_1} \dots x_{k_5}$ . Consider a partition of  $w$  into subwords  $w = a \cdot a_1 \cdot \dots \cdot a_s$ , where  $s = \overline{5, 10}$ ,  $a_1, \dots, a_s \in F_d^\#$ ,  $a \in F_d$ , which is the result of intersection of partitions  $w = \bar{t}_1 = c_1 \cdot \psi_1(x_{j_1}) \cdot \dots \cdot \psi_1(x_{j_5})$ ,  $w = \bar{t}_2 = c_2 \cdot \psi_2(x_{k_1}) \cdot \dots \cdot \psi_2(x_{k_5})$ . Here we assume  $a \neq 1$  if and only if  $c_1, c_2$  are non-empty words. Then  $c_i = ad_i$ , where  $d_i \in F_d$ ,  $i = 1, 2$ . There is a permutation  $\sigma \in S_s$  such that if  $x_i < x_j$ , then  $a_{\sigma(i)} < a_{\sigma(j)}$ , where  $i, j = \overline{1, s}$ . Since  $d_1 \psi_1(t''_1), d_2 \psi_2(t''_2)$  are elements generated by words  $a_1, \dots, a_s$ , the substitutions  $d_i \psi_i(t''_i)|_{a_{\sigma(j)} \rightarrow x_j, j=\overline{1, s}} = t'_i$ ,  $i = 1, 2$ , are well-defined. It is easy to see

that for  $t'_i$  there is  $\psi'_i \in \mathcal{M}_{5,s}$ ,  $e_i \in F_s$  such that  $e_i \psi'_i(t''_i) = t'_i$ , i.e.  $t'_i \in M_s$ , where  $i = 1, 2$ . Define  $\phi \in \mathcal{M}_{s,d}$  by the following way:  $\phi(x_j) = a_{\sigma(j)}$ ,  $j = \overline{1, s}$ . We have  $a\phi(t'_i) = t_i$ , where  $i = 1, 2$ , thus the claim is proven.  $\triangle$

**Proposition 3** *If  $B(M_d)$  is a basis for  $N_{3,d}(1^d)$  for  $d = \overline{5, 10}$ , then  $B(M_d)$  is a basis for  $N_{3,d}(1^d)$  for any  $d \geq 5$ .*

**Proof.** Let  $d \geq 11$ . Consider  $t_1, t_2 \in M_d$ , where  $\bar{t}_1 = \bar{t}_2$ . We apply Lemma 8 to  $t_1, t_2$ ; further we use the notation from Lemma 8. By the data and Lemma 7  $M_s$  is complete under composition, hence  $t'_1 - t'_2 = \sum_i \alpha_i g'_i$ , where  $\alpha_i \in K$ ,  $g'_i \in M_s$ ,  $\bar{t}'_1 > \bar{g}'_i$ . Thus  $t_1 - t_2 = \sum_i \alpha_i g_i$ , where  $g_i = a\phi(g'_i) \in M_d$ ; because of the composition of monotonous substitutional mappings is a monotonous substitutional mapping. By monotony of  $\phi$ , we have  $\bar{t}_1 > \bar{g}_i$  for all  $i$ . Hence  $M_d$  is complete under composition, and Lemma 7 concludes the proof.  $\triangle$

## 7 Multilinear homogeneous component

**Notation.** For  $p = 2, 3$  and  $d \geq 1$  recursively define sets  $B_{1^d}$  of the words of multidegree  $1^d$ .

Let  $p = 2$ . Then

1)  $B'_1 = \{x_1\}$ ;

2) for  $d \geq 2$  define  $B'_{1^d} = x_1\{B_{1^{d-1}}|_{x_i \rightarrow x_{i+1}, i=\overline{1, d-1}}\} \cup \{\underline{e}_{d,k} | k = \overline{2, d}\} \cup \{\underline{f}_d\} \cup \{\underline{h}_{d,k} | k = \overline{3, d}\}$ .

Here, if  $i \geq 1, i \neq 4$ , then  $B_{1^i} = B'_{1^i}$ ;  $B_{1^4} = B'_{1^4} \cup \{x_2 x_1 x_4 x_3\}$  and

$\underline{e}_{d,k} = x_2 \cdots x_k \cdot x_1 \cdot x_{k+1} \cdots x_d$  ( $d \geq 2, k = \overline{2, d}$ );

$\underline{f}_d = x_2 \cdots x_{d-2} \cdot x_d x_1 x_{d-1}$  ( $d \geq 3$ );

$\underline{h}_{d,k} = x_k \cdot x_1 \cdots \tilde{x}_k \cdots x_d$  ( $d \geq 3, k = \overline{3, d}$ ).

Let  $p = 3$ . Then

1)  $B_1 = \{x_1\}$ ;

2) for  $d \geq 2$  define  $B_{1^d} = x_1\{B_{1^{d-1}}|_{x_i \rightarrow x_{i+1}, i=\overline{1, d-1}}\} \cup x_2\{B_{1^{d-1}}|_{x_i \rightarrow x_{i+1}, i=\overline{2, d-1}}\} \cup \{\underline{e}_{d,k} | k = \overline{3, d}\}$ .

Here,  $\underline{e}_{d,k} = x_3 \cdots x_k \cdot x_1 x_2 \cdot x_{k+1} \cdots x_d$  ( $d \geq 3, k = \overline{3, d}$ ).

For future needs define  $B_{1^0} = \{1\}$ , where 1 stands for the empty word.

The aim of this section is to prove the following theorem:

**Theorem 1** *For  $p = 2, 3, d \geq 1$  the set  $B_{1^d}$  is a basis for  $N_{3,d}(1^d)$ .*

**Remark 2** *It is not difficult to see that*

$$|B_{1^d}| = \begin{cases} d(d-1) & , \quad p = 2, d \geq 4 \\ 2^d - d & , \quad p = 3 \end{cases}.$$

Let  $V$  be a finite dimensional vector space over  $K$ ,  $V = \text{lin}\{v_1, \dots, v_m\}$ , where non-zero vectors  $\{v_i\}$  are linearly ordered by the following way:  $v_1 < \dots < v_m$ . Note that the vectors  $v_1, \dots, v_m$  need not be linearly independent.

**Definition.** A basis of  $V$   $v_{k_1}, \dots, v_{k_s}$  is called *minimal* (with respect to the linearly ordered set  $v_1, \dots, v_m$ ), if for any  $i = \overline{1, m}$  we have  $v_i = \sum_j \alpha_{ij} v_{k_j}$ ,  $\alpha_{ij} \in K$ , where  $k_j \leq i$ .

Consider  $L_1 = \{v_{j_1}, \dots, v_{j_s}\}$ ,  $L_2 = \{v_{k_1}, \dots, v_{k_s}\}$  which are bases for  $V$ , where  $j_1 < \dots < j_s$ ,  $k_1 < \dots < k_s$ . We write  $L_1 < L_2$ , if there is  $l = \overline{1, s}$  such that  $j_1 = k_1, \dots, j_{l-1} = k_{l-1}, j_l < k_l$ .

**Lemma 9** 1. *A basis of  $V$   $v_{k_1}, \dots, v_{k_s}$  is minimal (with respect to the linearly ordered set  $v_1, \dots, v_m$ ) if and only if it is the least one with respect to the determined linear order.*

2. *The minimal basis is uniquely determined.*

**Proof.** 1. Let  $L$  be the minimal basis. Then

$$\text{if } v_k \notin \text{lin}\{v_1, \dots, v_{k-1}\}, \text{ then } v_k \in L; \text{ else } v_k \notin L \text{ (} k = \overline{1, m} \text{)}. \quad (4)$$

Thus,  $L$  is the least basis.

Let  $L \subset \{v_1, \dots, v_m\}$  be the least basis. Then condition (4) is valid for it. Thus  $L$  is the minimal basis.

2. This item follows from item 1.  $\triangle$

Apply aforesaid on minimal bases to  $N_{3,d}(1^d)$ . As a linearly ordered set we take  $\{w \in F_d \mid \text{mdeg}(w) = 1^d\}$ .

**Lemma 10** 1.  *$B(M_5)$  is a basis for  $N_{3,d}(1^5)$ .*

2. *Let  $d \geq 5$ . If  $B(M_d)$  is a basis for  $N_{3,d}(1^d)$ , then  $B(M_d)$  is the minimal basis.*

**Proof.** 1. The definition of  $M_5$  implies that  $M_5$  is complete under composition. Lemma 7 concludes the proof.

2. Identities  $M_d$  imply that any word, of  $N_{3,d}$ , which do not belong to  $B(M_d)$  can be expressed in terms of lower words. Thus  $B(M_d)$  is the minimal basis.  $\triangle$

**Lemma 11** 1. If  $p = 2$ , then for  $i = 4, 5$  the minimal basis for  $N_{3,i}(1^i)$  is  $B_{1^i}$ , and

$$B_{1^4} = \left\{ \begin{array}{l} 1234, 1243, 1324, 1342, 1423, \\ 2134, 2143, 2314, 2341, 2413, \\ 3124, 4123 \end{array} \right\};$$

$$B_{1^5} = B(M_5) = \left\{ \begin{array}{l} 12345, 12354, 12435, 12453, 12534, \\ 13245, 13254, 13425, 13452, 13524, \\ 14235, 15234, \\ 21345, 23145, 23415, 23451, \\ 23514, \\ 31245, 41235, 51234 \end{array} \right\}.$$

2. If  $p = 3$ , then for  $i = 4, 5$  the minimal basis for  $N_{3,i}(1^i)$  is  $B_{1^i}$ , and

$$B_{1^4} = \left\{ \begin{array}{l} 1234, 1243, 1324, 1342, 1423, \\ 2134, 2143, 2314, 2341, 2413, \\ 3124, 3412 \end{array} \right\};$$

$$B_{1^5} = B(M_5) = \left\{ \begin{array}{l} 12345, 12354, 12435, 12453, 12534, \\ 13245, 13254, 13425, 13452, 13524, \\ 14235, 14523, \\ 21345, 21354, 21435, 21453, 21534, \\ 23145, 23154, 23415, 23451, 23514, \\ 24135, 24513, \\ 31245, 34125, 34512 \end{array} \right\}.$$

**Proof.** By Lemma 10,  $B(M_5)$  is the minimal basis for  $N_{3,5}(1^5)$ , and, in particular, it does not depend on the choice of  $M_5$  (see Lemma 9). Expressing higher words in terms of lower words by Gauss's method, we solve the system  $\mathcal{S}_{1^d}$  and find the minimal basis for  $N_{3,d}(1^d)$ . These calculations were performed by means of a computer programme for  $d = 4, 5$ ,  $p = 2, 3$ .  $\triangle$

**Lemma 12** If  $p = 2$ ,  $d \geq 5$ , then

$$\overline{M}_d = \left\{ \begin{array}{l} w_3 = ua_1a_2a_3, \quad a_1 > a_2 > a_3, \\ w_{4,1} = ua_1a_2a_3a_4, \quad a_1 > a_2, a_4 \text{ and } a_3 > a_4, \\ w_{4,2} = ua_1a_2a_3a_4, \quad a_1 > a_4 \text{ and } a_2 > a_3, \\ w_{5,1} = ua_1a_2a_3a_4a_5, \quad a_1 > a_2 \text{ and } a_3 > a_4 \text{ or } a_3 > a_5 \text{ or } a_4 > a_5, \\ w_{5,2} = ua_1a_2a_3a_4a_5, \quad a_1 > a_3 \text{ and } a_2 > a_4 \text{ or } a_2 > a_5 \text{ or } a_4 > a_5, \\ w_{5,3} = ua_1a_2a_3a_4a_5, \quad a_1 > a_4 \text{ and } a_2 > a_5. \end{array} \right\},$$

where all elements of  $\overline{M}_d$  are words of multidegree  $1^d$ ,  $u \in F_d$ ,  $a_i \in F_d^\#$ ,  $i = \overline{1, 5}$ .

If  $p = 3$ ,  $d \geq 5$ , then

$$\overline{M}_d = \left\{ \begin{array}{ll} w_3 = ua_1a_2a_3, & a_1 > a_2 > a_3, \\ w_4 = ua_1a_2a_3a_4, & a_1 > a_2, a_4, \\ w_{5,1} = ua_1a_2a_3a_4a_5, & a_1 > a_2, a_3 \text{ and } a_4 > a_5, \\ w_{5,2} = ua_1a_2a_3a_4a_5, & a_1 > a_3, a_4 \text{ and } a_2 > a_5, \\ w_{5,3} = ua_1a_2a_3a_4a_5, & a_1 > a_4, a_5 \text{ and } a_2 > a_3. \end{array} \right\},$$

where all elements of  $\overline{M}_d$  are words of multidegree  $1^d$ ,  $u \in F_d$ ,  $a_i \in F_d^\#$ ,  $i = \overline{1, 5}$ .

**Proof.** It is sufficient to prove the statement for  $d = 5$ . Denote by  $M$  the set from the formulation of the Lemma. Considering all possibilities, we get that  $B(M_5) = \{w \in F_5 \mid \text{mdeg}(w) = 1^5, w \notin M\}$  (see Lemma 11).  $\triangle$

**Lemma 13** For  $p = 2, 3$ ,  $d \geq 5$  we have  $B(M_d) = B_{1^d}$ .

**Proof.** For  $p = 2, 3$  we get  $B_{1^d} \subset B(M_d)$  by Lemma 12.

Further  $w_{i,j}, w_i$  stand for the words from Lemma 12.

**The case of  $p = 2$ .** Inclusion  $B(M_d) \subset B_{1^d}$  follows from items 1, 2 (see below) by induction on  $d$ .

1. If  $w = x_{i_1} \cdots x_{i_d} \in B(M_d)$ ,  $i_1 \geq 3$ , then  $w = \underline{h}_{d,i_1}$ .

Proof. Let  $w = x_{i_1}u$ . The word  $w$  contains letters  $x_1, x_2$ . Denote by  $u_1, u_2, u_3$  some elements of  $F_d$ . If  $u = u_1x_2u_2x_1u_3$ , then  $w = w_3 \in \overline{M}_d$ , that is a contradiction. Hence  $u = u_1x_1u_2x_2u_3$ . If the word  $u_1$  is not empty, then  $w = w_{4,2} \in \overline{M}_d$ , that is a contradiction. If the word  $u_2$  is not empty, then  $w = w_{4,1} \in \overline{M}_d$ , that is a contradiction. Assume that  $u_3 = x_{j_1} \cdots x_{j_s}$  and there are  $k, t$  such that  $k < t \leq s$  and  $j_k > j_t$ . Then  $w = w_{5,1} \in \overline{M}_d$ , that is a contradiction. Hence  $w = \underline{h}_{d,i_1}$ .

2. If  $w = x_2x_{i_2} \cdots x_{i_d} \in B(M_d)$ , then  $w = \underline{e}_{d,k}$  for some  $k = \overline{2, d}$  or  $w = \underline{f}_d$ .

Proof. Consider  $w = x_2u_1x_1u_2$ , where  $u_1, u_2 \in F_d$ . If there are not  $r, s$  such that  $r > s$  and  $u_1u_2 = v_1x_rv_2x_sv_3$ , where  $v_1, v_2, v_3 \in F_d$ , then  $w = \underline{e}_{d,k}$  for some  $k = \overline{2, d}$ . Assume that there are such  $r, s$ . If the word  $v_3$  contains the letter  $x_1$ , then  $w = w_{4,2} \in \overline{M}_d$ ; a contradiction. If the word  $v_1$  contains the letter  $x_1$ , then  $w = w_{5,1}$  or  $w = w_{5,2}$ , hence  $w \in \overline{M}_d$ ; a contradiction. Let the word  $v_2$  contains the letter  $x_1$ . If the word  $v_3$  is not empty, then  $w = w_{5,2} \in \overline{M}_d$ ; a contradiction. If  $\text{deg}(v_2) > 1$ , then  $w = w_{5,2}$  or  $w = w_{5,3}$ , hence  $w \in \overline{M}_d$ ; a contradiction. There is the only possibility which we have not considered, namely  $w = \underline{f}_d$ .

**The case of  $p = 3$ .** Inclusion  $B(M_d) \subset B_{1^d}$  follows from items 1, 2 (see below) by induction on  $d$ .

1. If  $w = x_{i_1} \cdots x_{i_d}$ ,  $i_1 \geq 4$ , then  $w \notin B(M_d)$ .

Proof. The word  $w_1 = x_{i_2} \cdots x_{i_d}$  contains the letters  $x_1, x_2, x_3$ . There are  $r, s = \overline{1, 3}$  such that the word  $w_1$  contains some letters between letters  $x_r, x_s$ . We also have  $x_{i_1} > x_r, x_s$ . Thus  $w = w_4 \in \overline{M}_d$ , that is  $w \notin B(M_d)$ .

2. If  $w = x_3x_{i_2} \cdots x_{i_d} \in B(M_d)$ , then  $w = \underline{e}_{d,k}$  for some  $k = \overline{3, d}$ .

Proof. If  $w = x_3u_1x_2u_2x_1u_3$  for some elements  $u_1, u_2, u_3$  of  $F_d$ , then  $w = w_3 \in \overline{M}_d$ ; a contradiction. Thus  $w = x_3u_1x_1u_2x_2u_3$ . If the word  $u_2$  is not empty, then  $w = w_4 \in \overline{M}_d$ , that is a contradiction. Hence  $u = x_3u_1x_1x_2u_3$ . If there are  $r, s$  such that  $r > s$  and  $u_1u_3 = v_1x_rv_2x_sv_3$  for some  $v_1, v_2, v_3 \in F_d$ , then  $w = w_{5,1}$  or  $w = w_{5,2}$  or  $w = w_{5,3}$ . Therefore  $w \in \overline{M}_d$ , that is a contradiction. Hence there are not such  $r, s$ . So  $w = \underline{e}_{d,k}$  for some  $k = \overline{3, d}$ .  $\triangle$

**Lemma 14** Let  $p = 3$ ,  $d \geq 6$ ,  $\phi_k$  is the mapping from Lemma 5, where  $k = \overline{1, d}$ . Then for any  $w \in B(M_d)$  we have  $\phi_k(w)|_{x_i \rightarrow x_{i-1}, i=\overline{k+1, d}} \in B(M_{d-1})$ .

**Proof.** If for a word  $w$  of multidegree  $1^d$  we have  $\phi_k(w)|_{x_i \rightarrow x_{i-1}, i=\overline{k, d}} \in \overline{M}_{d-1}$ , then  $w \in \overline{M}_d$ , that is  $w \notin B(M_d)$ .  $\triangle$

**Proof of theorem 1.** If  $d = 1, 2, 3$  then obviously  $B_{1^d}$  is a basis. For  $d = 4, 5$   $B_{1^d}$  is a basis by Lemma 11. Proposition 3 and Lemma 13 imply that in order to prove the Theorem it is sufficient to verify that  $B_{1^d}$  is linearly independent in  $N_{3,d}$  for  $d = \overline{6, 10}$ . This verification was done by means of a computer programme applying the algorithm described below.

**The case of  $p = 2$ .** Assume that there is an identity  $f = \sum_{w \in B_{1^d}} \alpha_w w$ ,  $\alpha_w \in K$ , such that  $f = 0$  in  $N_{3,d}$ . Considering  $\phi_{ij}(t) = 0$ ,  $\phi(t) = 0$ , where  $\phi_{ij}, \phi$  are mappings from Lemma 5, we get a homogeneous system of linear equations in  $\{\alpha_w\}$ . Having solved this system we get that  $\alpha_w = 0$  for any  $w \in B_{1^d}$ . It was calculated for  $d = \overline{6, 10}$  by means of a computer programme.

**The case of  $p = 3$ .** By Lemma 13, we have  $B_{1^d} = B(M_d)$ . Identities from  $M_5$  express elements of the set  $\{w \in F_5 \mid \text{mdeg}(w) = 1^5\}$  in terms of elements of  $B_{1^5}$ . Applying these identities, we get that for any word  $w$  of multidegree  $1^d$   $w = f_w$  in  $N_{3,d}$ , where  $f_w \in \text{lin } B_{1^d}$ . Applying identities  $\{w = f_w\}$ , rewrite identities  $\{f_1T_3(a_1, a_2, a_3)f_2 \mid f_1, f_2 \in F_d, \deg(a_1) \leq 3, \deg(a_2) = \deg(a_3) = 1\}$  in terms of linear combinations of the elements of  $B_{1^d}$ . As a result, we get only trivial identities. This, together with Lemma 4, imply that the system of identities  $\mathcal{S}_{1^d}$  is equivalent to the set of identities  $M = \{w - f_w \mid w \in F_d, \text{mdeg}(w) = 1^d, w \notin B_{1^d}\}$ . The set  $M$  is linearly independent in  $K\langle F_d \rangle$ , thus  $B_{1^d}$  is linearly independent in  $N_{3,d}$ . The given algorithm performed by means of a computer programme proved that  $B_{1^d}$  is linearly independent in  $N_{3,d}$  when  $d = \overline{6, 9}$ . Here we need Lemma 4 in order to decrease the quantity of identities which have to be considered. For  $d = 10$  described algorithm ran for a long time, thus we used another approach to the case of  $d = 10$ .

Assume that there is an identity  $f = \sum_{w \in B_{1^d}} \alpha_w w$ , where  $\alpha_w \in K$ , such that  $f = 0$  in  $N_{3,d}$ . Consider mappings  $\phi_k$  ( $k = \overline{1, d}$ ),  $\phi_+$  from Lemma 5. We assume



that  $B_{1^{d-1}}$  is linearly independent in  $N_{3,d}$ . Hence, we get a homogeneous system of linear equations in  $\{\alpha_w\}$  (see also Lemma 14). For even  $d = \overline{6, 10}$  it was calculated by means of a computer programme that this system has the only solution  $\alpha_w = 0$  for any  $w \in B_{1^d}$ .  $\triangle$

## 8 The case of $p = 2$

Let  $d \geq 4$ ,  $i, j = \overline{2, d}$ ,  $i \neq j$ . Introduce notations for some words of multidegree  $21^{d-1}$ :  $\underline{a}_i = x_1^2 x_i x_2 \cdots \tilde{x}_i \cdots x_d$ ,  $\underline{b}_i = x_1 \cdots \tilde{x}_i \cdots x_d x_i x_1^2$ ,  $\underline{c}_{ij} = x_i x_1^2 x_j x_2 \cdots \tilde{x}_i \cdots \tilde{x}_j \cdots x_d$ , if  $i < j$ , and  $\underline{c}_{ij} = x_i x_1^2 x_j x_2 \cdots \tilde{x}_j \cdots \tilde{x}_i \cdots x_d$ , if  $i > j$ . By items 1, 7 of Lemma 1 we have

$$N_{3,d}(21^{d-1}) = \text{lin}\{\underline{a}_i, \underline{b}_i, \underline{c}_{ij} \mid i, j = \overline{2, d}, i \neq j\}. \quad (5)$$

**Theorem 2** *Let  $p = 2$ ,  $d \geq 1$ .*

1. *For  $\Delta = (\delta_1, \dots, \delta_d)$ ,  $d \leq 3$  a basis for  $N_{3,d}(\Delta)$  is the set  $B_\Delta$  defined in Proposition 1.*
2. *For  $d \geq 4$  a basis for  $N_{3,d}(1^d)$  is the set  $B_{1^d}$  defined above.*
3. *A basis for  $N_{3,4}(21^3)$  is the set  $B_{21^3} = \{\underline{a}_i, \underline{b}_i, \underline{c}_{23}, \underline{c}_{32} \mid i = \overline{2, 4}\}$ .*  
*For  $d \geq 5$  a basis for  $N_{3,d}(21^{d-1})$  is the set  $B_{21^{d-1}} = \{\underline{a}_i, \underline{b}_i, \underline{c}_{23} \mid i = \overline{2, d}\}$ .*
4. *For  $d \geq 4$  a basis for  $N_{3,d}(2^2 1^{d-2})$  is the set  $B_{2^2 1^{d-2}} = \{x_1^2 x_2^2 x_3 \cdots x_d, x_2^2 x_1^2 x_3 \cdots x_d\}$ .*
5. *For  $d \geq 4$  a basis for  $N_{3,d}(31^{d-1})$  is the set  $B_{31^{d-1}} = \{x_1^2 x_2 \cdots x_d x_1\}$ .*
6. *The rest of  $\mathcal{N}^d$ -homogeneous components of  $N_{3,d}$  are equal to zero.*

In order to prove item 3 we need the following Lemma.

Denote  $h_{ij} = \underline{b}_i + \underline{c}_{ij} + \underline{a}_j$ . Consider identities of multidegree  $21^{d-1}$ :

$$\begin{aligned} M_0 &= \{f_1 I_i(x_1, a, b, c) f_2 \mid i = \overline{1, 3}\}, \\ M_1 &= \text{the set of identities (a) from Lemma 2}, \\ M_2 &= \{f_1 T_3(x_1^2, a, b) f_2\}, \\ M_3 &= \{T_3(x_1^2, x_i, x_j) a, a T_3(x_1^2, x_i, x_j) \mid 2 \leq i < j \leq d\}, \\ M_4 &= \begin{cases} \{h_{23} + h_{34}, h_{23} + h_{42}, h_{32} + h_{24}, h_{32} + h_{43}\}, & \text{if } d = 4 \\ \{h_{23} + h_{ij} \mid 2 \leq i \neq j \leq d, i \neq 2 \text{ or } j \neq 3\}, & \text{if } d \geq 5, \end{cases} \end{aligned}$$

where  $a, b, c \in F_d^\#$ ,  $f_1, f_2 \in F_d$ . For  $i = \overline{1, 4}$  denote  $L_i = \text{lin}\{M_0 \cup M_i\}$ .

**Lemma 15** *Let  $p = 2$ ,  $d \geq 4$ . Then*

1.  $L_1 = L_2$ .

2.  $L_2 = L_3$ .

3.  $L_3 = L_4$ .

**Proof.** Let  $i, j, k, l \in \overline{2, d}$  be pairwise different numbers.

1. Inclusion  $L_2 \subset L_1$  is obvious.

Consider an identity  $t \in M_1$ .

If  $t = T_3(ax_1^2b, c, d)$ , then  $t = aI_1(x_1, b, c, d) + I_2(x_1, c, d, a)b + I_3(x_1, ca, b, d) + dI_3(x_1, a, b, c) + 2dcax_1^2b$ .

If  $t = T_3(x_1^2a, b, c)$ , then  $t = I_1(x_1, a, b, c) + I_3(x_1, c, a, b) + I_3(x_1, b, a, c)$ .

If  $t = x_1^2aT_3(b, c, d)$ , then  $t = I_1(x_1, a, b, c)d + I_1(x_1, a, b, d)c + I_1(x_1, a, c, d)b$ .

If  $t = x_1^2T_3(a, b, c)$ , then  $t = I_1(x_1, a, b, c) + I_1(x_1, b, a, c) + I_1(x_1, c, a, b)$ .

Thus we get  $L_1 \subset L_2$ .

3. Introduce notations for identities:

$$\begin{aligned} f_{ijk} &= \underline{a}_i + \underline{a}_j + \underline{c}_{ij} + \underline{c}_{ji} + \underline{c}_{ik} + \underline{c}_{jk} = T_3(x_1^2, x_i, x_j)x_k a + \{M_0\}, \\ g_{ijk} &= \underline{b}_j + \underline{b}_k + \underline{c}_{jk} + \underline{c}_{kj} + \underline{c}_{ij} + \underline{c}_{ik} = ax_i T_3(x_1^2, x_j, x_k) + \{M_0\}, \end{aligned}$$

where  $a \in F_d$ . We have

$$f_{ijk} + g_{jik} = h_{ki} + h_{ij} \in L_3. \quad (6)$$

If  $d = 4$ , then formula (6) implies  $L_4 \subset L_3$ .

If  $d \geq 5$ , then replacing indices in formula (6) we get  $h_{ij} + h_{jk} \in L_3$ ,  $h_{ij} + h_{jl} \in L_3$ ,  $h_{jl} + h_{lk} \in L_3$ ,  $h_{lk} + h_{kj} \in L_3$ . Add up last four formulas and get  $h_{jk} + h_{kj} \in L_3$ . Thus  $h_{ij} + h_{ji}, h_{ij} + h_{jk} \in L_3$ . Therefore  $L_4 \subset L_3$ .

From  $f_{ijk} = 4h_{23} + h_{ij} + h_{ji} + h_{ik} + h_{jk} \in L_4$ ,  $g_{ijk} = 4h_{23} + h_{jk} + h_{kj} + h_{ik} + h_{ij} \in L_4$  we can see that  $L_3 \subset L_4$  for  $d \geq 5$ .

Let  $d = 4$ . Equalities  $f_{ijk} = f_{jik}$ ,  $g_{ijk} = g_{ikj}$ ,  $f_{234} + f_{243} = f_{342}$ ,  $g_{234} + g_{324} = g_{423}$ ,

$$\begin{aligned} f_{234} &= (h_{23} + h_{34}) + (h_{32} + h_{24}), \\ f_{243} &= (h_{23} + h_{42}) + (h_{32} + h_{24}) + (h_{32} + h_{43}), \\ g_{234} &= (h_{23} + h_{34}) + (h_{32} + h_{24}) + (h_{32} + h_{43}), \\ g_{324} &= (h_{23} + h_{34}) + (h_{23} + h_{42}) + (h_{32} + h_{24}), \end{aligned}$$

imply  $L_3 \subset L_4$ .

2. Inclusion  $L_3 \subset L_2$  is obvious.

Consider an identity  $t \in M_2$ .

If  $t = T_3(x_1^2, a, bc)$ , then  $t = T_3(x_1^2, a, b)c + bT_3(x_1^2, a, c) + \{M_0\}$ .

If  $d \geq 5$ , then for  $t = f_1x_kT_3(x_1^2, x_i, x_j)x_l f_2$ ,  $f_1, f_2 \in F_d$ , we have  $t = \underline{c}_{ki} + \underline{c}_{kj} + \underline{c}_{ij} + \underline{c}_{ji} + \underline{c}_{il} + \underline{c}_{jl} + \{M_0\} \in L_4$ ; thus  $t \in L_3$  by item 3.

Therefore  $L_2 \subset L_3$ .  $\triangle$

**Proof of theorem 2.** 2. See Theorem 1.

3. By equality (5) and identities  $M_4$  it is sufficient to show that  $B_{21^{d-1}}$  is linear independent in  $N_{3,d}$ .

Lemmas 2 and 15 imply that all identities, of  $N_{3,d}(21^{d-1})$ , generated by  $x_1^2, x_2, \dots, x_d$  are consequences of identities  $M_0 \cup M_4$ . Let  $f$  be a mapping such that the image of a word  $w$ , generated by  $x_1^2, x_2, \dots, x_d$ , of multidegree  $21^{d-1}$  is equal to the result of application of the identities from item 7 of Lemma 1 to  $w$ , i.e.  $f(w)$  is equal to  $\underline{a}_i$ ,  $\underline{b}_i$  or  $\underline{c}_{ij}$  for some  $i, j$ . Identities  $M_0 \cup M_4$  are equivalent to the identities  $M'_0 \cup M_4 = M$ , where  $M'_0 = \{w + f(w) | w \text{ is a word, } \text{mdeg}(w) = 21^{d-1}, w \neq f(w)\}$ . Every identity of  $M'_0$  ( $M_4$ , respectively) contains a word which is not a summand of any element of  $B_{21^{d-1}}$  and is a summand of one and only one identity of  $M$  ( $M_4$ , respectively). Moreover we can assume that regarded words are pairwise different. Thus item 2 of Remark 1 implies that  $B_{21^{d-1}}$  is linearly independent in  $N_{3,d}$ .

4. Denote  $a = x_1^2 x_2^2 x_3 \cdots x_d$ ,  $b = x_2^2 x_1^2 x_3 \cdots x_d$ . By items 1, 4 and 7 of Lemma 1, we have  $\text{lin}\{a, b\} = N_{3,d}(2^2 1^{d-2})$ . We claim that  $a, b$  are linearly independent in  $N_{3,d}$ . Consider the homomorphism of vector spaces  $\psi : K\langle F_d \rangle \rightarrow K\langle F_d \rangle(2^2 1^{d-2})$  defined by the following way: for a word  $w$   $\psi(w) = \alpha a + \beta b$ , where  $\alpha$  ( $\beta$ , respectively) is equal to the number of subwords  $x_1 x_2$  ( $x_2 x_1$ , respectively) in the word  $w$ .

For  $u, v \in F_d^\#$  define

$$\psi(u, v) = \begin{cases} a & , \text{ if } u = u_1 x_1, v = x_2 v_1 (u_1, v_1 \in F_d) \\ b & , \text{ if } u = u_1 x_2, v = x_1 v_1 (u_1, v_1 \in F_d) \\ 0 & , \text{ otherwise} \end{cases}.$$

It is easy to see that for  $u_1, \dots, u_s \in F_d^\#$

$$\psi(u_1 \cdots u_s) = \sum_{i=1}^s \psi(u_i) + \sum_{i=1}^{s-1} \psi(u_i, u_{i+1}). \quad (7)$$

Consider an identity  $t$  of  $\mathcal{S}_{2^2 1^{d-2}}$ .

If  $t = f_1 T_1(g) f_2$ ,  $f_1, f_2 \in F_d$ ,  $g \in F_d^\#$ , then  $t \notin \mathcal{S}_{2^2 1^{d-2}}$ . It is a contradiction.

If  $t = f_1 T_2(g_1, g_2) f_2$ ,  $f_1, f_2 \in F_d$ ,  $g_1, g_2 \in F_d^\#$ , then  $\psi(t) = \psi(g_2) + \psi(f_1) + \psi(f_2) + \psi(f_1, g_2) + \psi(g_2, f_2)$ , by equality (7). The multidegree of  $t$  is  $2^2 1^{d-2}$ , thus  $g_1 \in \{x_1, x_2, x_1 x_2, x_2 x_1\}$ . Hence  $\psi(t) = 0$ .

If  $t = f_1 T_3(g_1, g_2, g_3) f_2$ ,  $f_1, f_2 \in F_d$ ,  $g_1, g_2, g_3 \in F_d^\#$ , then  $\psi(t) = 0$  by equality (7).

Therefore  $\psi(t) = 0$  for any identity  $t$  of  $N_{3,d}(2^2 1^{d-2})$ .

If  $t = \alpha a + \beta b$ ,  $\alpha, \beta \in K$ , is an identity of  $N_{3,d}$ , then  $\psi(t) = \alpha a + \beta b = 0$  in  $K\langle F_d \rangle$ . Hence  $\alpha = \beta = 0$ .

5. By Lemma 1 we have  $u = x_1^2 x_2 \cdots x_d x_1 \neq 0$ . Identities from items 1, 7 of Lemma 1 imply that  $x_1^2 abcx_1 = x_1(abcx_1^2) = x_1 bacx_1^2 = x_1^2 bacx_1$  in  $N_{3,d}$ . The last identity together with item 7 of Lemma 1 imply that  $\text{lin}\{u\} = N_{3,d}(31^{d-1})$ .

6. It follows from item 3 of Lemma 1.  $\triangle$

## 9 The case of $p = 3$

**Notation.** For  $p = 3$ ,  $r, s, l \geq 0$  determine the set  $B_{3^r 1^s}$  of words of multidegree  $3^r 1^s$  and the set  $B_{3^r 2^s 1^l}$  of words of multidegree  $3^r 2^s 1^l$ :

$$\begin{aligned} B_{3^r 1^s} &= \underline{u}_{2r} \{B_{1^s} |_{x_i \rightarrow x_{i+2r}, i=\overline{1,s}}\} \cup \{\underline{q}_{2r,s,k} | k = \overline{1,s}\} \quad (r \geq 0), \\ B_{3^{2r+1} 1^s} &= \underline{u}_{2r} x_{2r+1}^2 x_{2r+2} \{B_{1^s} |_{x_1 \rightarrow x_{2r+1}, x_i \rightarrow x_{i+2r+1}, i=\overline{2,s}}\} \cup \{\underline{q}_{2r+1,s,k} | k = \overline{3, s+1}\} \cup \\ &\quad \{\underline{q}_{2r+1,s}\} \quad (r \geq 0), \text{ where} \\ \underline{q}_{2r,s,k} &= \underline{u}_{2r-2} \cdot x_{2r-1}^2 x_{2r}^2 \cdot x_{2r+1} \cdots x_{2r+k} \cdot x_{2r-1} x_{2r} \cdot x_{2r+k+1} \cdots x_{2r+s} \quad (r, s \geq 1, \\ &\quad k = \overline{1,s}), \\ \underline{q}_{2r+1,s,k} &= \underline{u}_{2r} \cdot x_{2r+1}^2 \cdot x_{2r+3} \cdots x_{2r+k} \cdot x_{2r+1} x_{2r+2} \cdot x_{2r+k+1} \cdots x_{2r+s+1} \quad (r \geq 0, s \geq 2, \\ &\quad k = \overline{3, s+1}), \\ \underline{q}_{2r+1,s} &= \underline{u}_{2r-2} \cdot x_{2r-1}^2 x_{2r}^2 x_{2r+1}^2 x_{2r-1} x_{2r} x_{2r+1} \cdot x_{2r+2} \cdots x_{2r+s+1} \quad (r \geq 1, s \geq 0), \\ \underline{u}_{2k} &= \underline{u}_{12} \cdots \underline{u}_{2k-1,2k} \quad (k \geq 1), \underline{u}_0 \text{ is the empty word,} \\ \underline{u}_{ij} &= x_i^2 x_j^2 x_i x_j. \end{aligned}$$

Define  $B_{3^r 2^s 1^l} = B_{3^r 1^{s+l}} |_{x_i \rightarrow x_i^2, i=\overline{r+1, r+s}}$ .

As an example we point out that  $B_{3^0 1^s} = B_{1^s}$ ,  $B_{3^{2r}} = \{\underline{u}_{2r}\}$ ,  $B_{3^{2r+1}} = \{\underline{q}_{2r+1,0}\}$ ,  $B_{3^{2r+1} 1} = \{\underline{u}_{2r} x_{2r+1}^2 x_{2r+2} x_{2r+1}, \underline{q}_{2r+1,1}\}$ .

**Theorem 3** *Let  $p = 3$ .*

1. *A basis for  $N_{3,d}(3^r 2^s 1^l)$  is the set  $B_{3^r 2^s 1^l}$ , where  $r, s, l \geq 0$ .*
2. *The rest of  $\mathcal{N}^d$ -homogeneous components of  $N_{3,d}$  are equal to zero.*

**Remark 3** *For  $p = 3$ ,  $r \geq 2$ ,  $s, l \geq 0$  we have  $|B_{3^r 2^s 1^l}| = 2^{s+l}$ .*

**Proof of theorem 3.** 1. Consider the homomorphism  $\phi : N_{3,d}(3^r 1^{s+l}) \rightarrow N_{3,d}(3^r 2^s 1^l)$ , defined by

$$\phi(x_i) = \begin{cases} x_i^2 & , \quad r+1 \leq i \leq r+s \\ x_i & , \quad \text{otherwise} \end{cases}.$$

By item 1 of Lemma 1  $\phi$  is surjective. By Lemma 3  $\phi$  is injective. Thus,  $\phi$  is an isomorphism of vector spaces. Hence it is sufficient to prove the Theorem for multidegree  $3^r 1^s$ . The last follows from Lemmas 16, 17 (see below).

2. See item 1 of Lemma 1.  $\triangle$

Further for multilinear elements  $f_1, f_2 \in K\langle F_d \rangle$ , where  $\deg(f_1 f_2) = m$ , writing  $f_1 \xi(f_2)$  means that  $\xi$  is a substitutional mapping of  $\mathcal{M}_{d,m}$  such that the multidegree of  $f_1 \xi(f_2)$  is equal to  $1^m$ .

Consider  $i_1, \dots, i_r \in \{1, 2\}$ , where  $r < d$ , and a word  $u$  such that

$$\begin{aligned} u &= x_1 & , \quad \text{if } r = d - 1 \\ u &\in \{\underline{e}_{d-r,k} \mid k = \overline{3, d-r}\} & , \quad \text{if } r < d - 1 \end{aligned}$$

Denote by  $(i_1 i_2 \dots i_r; u)$  the word  $w \in B_{1^d}$  which is the result of the following procedure. Let  $w_{r+1} = u \in B_{1^{d-r}}$ ,  $w_k = x_{i_k} \xi(w_{k+1}) \in B_{1^{d-k+1}}$  for every  $k = \overline{1, r}$ . Put  $w = w_1$ . For short, we will write  $(1^{s-1} i_s \dots i_r; u)$  instead of  $(1 \dots 1 i_s \dots i_r; u)$  and so on.

**Lemma 16** *Let  $r, s \geq 0$ . Then  $\text{lin } B_{3^r 1^s} = N_{3,d}(3^r 1^s)$ .*

**Proof.** Let  $J_{r,s} = M_{2r+s} \cup \{ux_{2i-1}x_{2i}v, ux_{2i}x_{2i-1}v, u(x_{2i}ax_{2i-1} + x_{2i-1}ax_{2i})v, u(x_{2i-1}x_{2j-1}x_{2i}bx_{2j} - x_{2i-1}x_{2j-1}x_{2i}x_{2j}b)v \mid u, v \in F_{2r+s}, a, b \in F_{2r+s}^\#, i, j = \overline{1, r}, b > x_{2j}\}$  be the subset of  $K\langle F_{2r+s} \rangle(1^{2r+s})$ .

Consider the homomorphism  $\phi : K\langle F_{2r+s} \rangle(1^{2r+s}) \rightarrow N_{3,d}(3^r 1^s)$ , defined by  $\phi(x_{2i-1}) = x_i^2$ ,  $\phi(x_{2i}) = x_i$ ,  $\phi(x_j) = x_{j-r}$ , where  $i = \overline{1, r}$ ,  $j = \overline{2r+1, 2r+s}$ . By item 1 of Lemma 1,  $\phi$  is surjective. Identities  $x^2 y^2 x a y = x^2 y^2 x y a$ ,  $x a x^2 + x^2 a x = 0$  of  $N_{3,d}$  (see Lemma 1) imply that  $\phi$  induces the epimorphism  $\phi_1 : K\langle F_{2r+s} \rangle(1^{2r+s}) / \text{lin}(J_{r,s}) \rightarrow N_{3,d}(3^r 1^s)$ . By equality (3) we have  $N_{3,d}(3^r 1^s) = \text{lin } \phi_1(B(J_{r,s}))$ . So in order to prove the statement it is sufficient to prove that

$$\phi_1(B(J_{r,s})) = B_{3^r 1^s}. \quad (8)$$

Note that  $B(J_{r,s}) = B(J_{r-1,s+2}) \setminus \overline{J}_{r,s}$  for  $r \geq 1$ ,  $s \geq 0$ .

For  $r = 0$  equality (8) is obvious.

Let  $r = 1$ . We have  $B(J_{1,s}) = B_{1^{s+2}} \setminus \overline{J}_{1,s} = x_1 \xi(B_{1^{s+1}}) \cup x_2 \xi(B_{1^{s+1}}) \cup \{x_3 \dots x_k \cdot x_1 x_2 \cdot x_{k+1} \dots x_{s+2} \mid k = \overline{3, s+2}\} \setminus \overline{J}_{1,s} = x_1 \xi(B_{1^{s+1}}) \setminus \overline{J}_{1,s} = x_1 x_2 \xi(B_{1^s}) \cup x_1 x_3 \xi(B_{1^s}) \cup \{x_1 \cdot x_4 \dots x_k \cdot x_2 x_3 \cdot x_{k+1} \dots x_{s+2} \mid k = \overline{4, s+2}\} \setminus \overline{J}_{1,s} = x_1 x_3 \xi(B_{1^s}) \cup \{x_1 \cdot x_4 \dots x_k \cdot x_2 x_3 \cdot x_{k+1} \dots x_{s+2} \mid k = \overline{4, s+2}\}$ . Hence equality (8) holds.

Let  $r = 2$ . We have  $B(J_{2,s}) = B(J_{1,s+2}) \setminus \overline{J}_{2,s} = \text{see above} = x_1 x_3 \xi(B_{1^{s+2}}) \cup \{x_1 \cdot x_4 \dots x_k \cdot x_2 x_3 \cdot x_{k+1} \dots x_{s+4} \mid k = \overline{4, s+4}\} \setminus \overline{J}_{2,s} = x_1 x_3 \xi(B_{1^{s+2}}) \setminus \overline{J}_{2,s} = x_1 x_3 x_2 \xi(B_{1^{s+1}}) \cup x_1 x_3 x_4 \xi(B_{1^{s+1}}) \cup \{x_1 x_3 \cdot x_5 \dots x_k \cdot x_2 x_4 \cdot x_{k+1} \dots x_{s+4} \mid k = \overline{5, s+4}\} \setminus \overline{J}_{2,s} = x_1 x_3 x_2 x_4 \xi(B_{1^s}) \cup x_1 x_3 x_2 x_5 \xi(B_{1^s}) \cup \{x_1 x_3 x_2 \cdot x_6 \dots x_k \cdot x_4 x_5 \cdot x_{k+1} \dots x_{s+4} \mid k =$

$\overline{6, s+4}\} \cup \{x_1x_3 \cdot x_5 \cdots x_k \cdot x_2x_4 \cdot x_{k+1} \cdots x_{s+4} \mid k = \overline{5, s+4}\} \setminus \overline{J_{2,s}} = x_1x_3x_2x_4\xi(B_{1s}) \cup \{x_1x_3 \cdot x_5 \cdots x_k \cdot x_2x_4 \cdot x_{k+1} \cdots x_{s+4} \mid k = \overline{5, s+4}\}$ . Thus equality (8) holds.

Let  $r = 3$ . We have  $B(J_{3,s}) = B(J_{2,s+2}) \setminus \overline{J_{3,s}} = \text{/see above/} = x_1x_3x_2x_4\xi(B_{1s+2}) \cup \{x_1x_3 \cdot x_5 \cdots x_k \cdot x_2x_4 \cdot x_{k+1} \cdots x_{s+6} \mid k = \overline{5, s+6}\} \setminus \overline{J_{3,s}} = x_1x_3x_2x_4\xi(B(J_{1,s})) \cup \{x_1x_3 \cdot x_5 \cdot x_2x_4 \cdot x_6 \cdots x_{s+6}\}$ . Thus equality (8) is proved to be true.

Let  $r \geq 4$ . By induction on  $r$  prove that

$$B(J_{r,s}) = x_1x_3x_2x_4\xi(B(J_{r-2,s})), \text{ where } r \geq 4. \quad (9)$$

Induction base. Let  $r = 4$ . We have  $B(J_{4,s}) = B(J_{3,s+2}) \setminus \overline{J_{4,s}} = \text{/see above/} = x_1x_3x_2x_4\xi(B(J_{1,s+2})) \cup \{x_1x_3x_5x_2x_4x_6 \cdot x_7 \cdots x_{s+8}\} \setminus \overline{J_{4,s}} = x_1x_3x_2x_4\xi(B(J_{2,s}))$ .

Induction step. We have  $B(J_{r,s}) = B(J_{r-1,s+2}) \setminus \overline{J_{r,s}} = \text{/induction hypothesis/} = x_1x_3x_2x_4\xi(B(J_{r-3,s+2})) \setminus \overline{J_{r,s}} = x_1x_3x_2x_4\xi(B(J_{r-2,s}))$ .

Formula (9) implies that equality (8) is valid.  $\triangle$

**Lemma 17** *Let  $r, s \geq 0$ . Then*

1. *The set  $B_{3^{2r+1}s}$  is linearly independent in  $N_{3,d}$ , where  $d = 2r + s$ .*
2. *The set  $B_{3^{2r+1}s}$  is linearly independent in  $N_{3,d}$ , where  $d = 2r + s + 1$ .*

**Proof.** Denote by  $\pi_i$  the homomorphism from Lemma 5. Note that for any  $k \geq 1$ ,  $v \in B_{1^k}$  words  $x_1x_2\xi(v)$ ,  $x_2x_1\xi(v)$  belong to  $B_{1^{k+2}}$ . Denote  $\pi_1 \cdots \pi_{2r-2}(\underline{u}_{2r-2}) = u$ .

1. If  $r = 0$ , then see Theorem 1. Let  $r \geq 1$ . For  $v \in B_{1^s}$ ,  $k = \overline{1, s}$  we have  $\pi_1 \cdots \pi_{2r}(\underline{u}_{2r}\xi(v)) = u(x_{2r-1}x_{2r} - x_{2r}x_{2r-1})\xi(v)$ ,  $\pi_1 \cdots \pi_{2r}(\underline{q}_{2r,s,k}) = u\xi(a - b - c + d)$ , where  $a = x_3 \cdots x_{k+2} \cdot x_1x_2 \cdot x_{k+3} \cdots x_{s+2} = \underline{e}_{s+2,k+2} \in B_{1^{s+2}}$ ,  $b = x_1 \cdot x_3 \cdots x_{k+2} \cdot x_2 \cdot x_{k+3} \cdots x_{s+2} = (12^k 1^{s-k}; x_1) \in B_{1^{s+2}}$ ,  $c = x_2 \cdot x_3 \cdots x_{k+2} \cdot x_1 \cdot x_{k+3} \cdots x_{s+2} = (2^{k+1} 1^{s-k}; x_1) \in B_{1^{s+2}}$ ,  $d = x_1x_2 \cdot x_3 \cdots x_{k+2} \cdot x_{k+3} \cdots x_{s+2} = (1^{s+1}; x_1) \in B_{1^{s+2}}$ . By above remark, for any  $w \in B_{3^{2r+1}s}$  we have  $\pi_1 \cdots \pi_{2r}(w) = \sum_i \alpha_{w,i} a_{w,i}$ , where  $\alpha_{w,i} \in K$ ,  $a_{w,i} \in B_{1^{2r+s}}$ . By item 1 of Remark 1, we get that the set  $\{\sum_i \alpha_{w,i} a_{w,i} \mid w \in B_{3^{2r+1}s}\}$  is linearly independent in  $K\langle F_d \rangle$ . So the assumption that  $B_{3^{2r+1}s}$  is linearly dependent in  $N_{3,d}$  gives that  $B_{1^{2r+s}}$  is linearly dependent in  $N_{3,d}$  (see Lemma 5), and the last contradicts Theorem 1.

2. Let  $v \in B_{1^s}$ ,  $k = \overline{3, s+1}$ . Denote  $a_v = \pi_1 \cdots \pi_{2r+1}(\underline{u}_{2r}x_{2r+1}^2x_{2r+2}\xi(v))$ ,  $b_k = \pi_1 \cdots \pi_{2r+1}(\underline{q}_{2r+1,s,k})$ ,  $c = \pi_1 \cdots \pi_{2r+1}(\underline{q}_{2r+1,s})$ . Let  $\phi_1$  be the homomorphism from item 4 of Lemma 5. We have that  $a_v = u(x_{2r-1}x_{2r} - x_{2r}x_{2r-1})(x_{2r+2}\xi(v) - x_{2r+1}x_{2r+2}\xi(\phi_1(v)))$ ,  $b_k = u(x_{2r-1}x_{2r} - x_{2r}x_{2r-1})(x_{2r+3} \cdots x_{2r+k} \cdot x_{2r+1}x_{2r+2} \cdot x_{2r+k+1} \cdots x_{2r+s+1} - x_{2r+1} \cdot x_{2r+3} \cdots x_{2r+k} \cdot x_{2r+2} \cdot x_{2r+k+1} \cdots x_{2r+s+1}) = u(x_{2r-1}x_{2r} - x_{2r}x_{2r-1})\xi(\underline{e}_{s+1,k} - (12^{k-2} 1^{s-k+1}; x_1))$ ,  $c = u(x_{2r-1}x_{2r+1}x_{2r} - x_{2r}x_{2r-1}x_{2r+1} + x_{2r}x_{2r+1}x_{2r-1} - x_{2r+1}x_{2r-1}x_{2r})x_{2r+2} \cdots x_{2r+s+1} = u\xi((121^s; x_1) - (21^{s+1}; x_1) +$

$(2^2 1^s; x_1) - \underline{e}_{s+3,3}$ ). Above remark and Lemma 14 imply that  $a_v, b_k, c \in \text{lin } B_{1^{2r+s+1}}$ . Thus  $a_v = \sum_i \alpha_{v,i} a_{v,i}$ ,  $b_k = \sum_i \beta_{k,i} b_{k,i}$ ,  $c = \sum_i \gamma_i c_i$ , where  $\alpha_{v,i}, \beta_{k,i}, \gamma_i \in K$ ,  $a_{v,i}, b_{k,i}, c_i \in B_{1^{2r+s+1}}$ . The set  $\{a_v, b_k, c | v \in B_{1^s}, k \in \overline{3, s+1}\}$  is linearly independent, because of each set from the class of sets  $\{\{a_{v,i}\}, \{b_{k,i}\}, \{c_i\} | v \in B_{1^s}, k \in \overline{3, s+1}\}$  contains an element which does not belong to other sets. Thus the assumption that  $B_{3^{2r+1} 1^s}$  is linearly dependent in  $N_{3,d}$  gives that  $B_{1^{2r+s+1}}$  is linearly dependent in  $N_{3,d}$  (see Lemma 5), and the last contradicts Theorem 1.  $\triangle$

## 10 Matrix invariants

Let  $n \geq 2$ . Denote by  $M_{n,d}(K) = M_n(K) \oplus \cdots \oplus M_n(K)$  the sum of  $d$  copies of the space of  $n \times n$  matrices. The general linear group  $GL_n(K)$  acts on  $M_{n,d}(K)$  by diagonal conjugation: for  $g \in GL_n(K)$ ,  $A_i \in M_n(K)$  ( $i = \overline{1, d}$ ) we have  $g(A_1, \dots, A_d) = (gA_1g^{-1}, \dots, gA_dg^{-1})$ . The coordinate ring of the affine space  $M_{n,d}(K)$  is the polynomial algebra  $K_{n,d} = K[x_{ij}(r) | 1 \leq i, j \leq n, r = \overline{1, d}]$ , where  $x_{ij}(r)$  stands for the function such that the image of  $(A_1, \dots, A_d) \in M_{n,d}(K)$  is  $(i, j)$ th entry of the matrix  $A_r$ . The action of  $GL_n(K)$  on  $M_{n,d}(K)$  induces the action on  $K_{n,d}$ :  $(g \cdot f)(A) = f(g^{-1}A)$ , where  $g \in GL_n(K)$ ,  $f \in K_{n,d}$ ,  $A \in M_{n,d}$ . Denote by  $R_{n,d} = \{f \in K_{n,d} | \text{for all } g \in GL_n(K) : gf = f\}$  the matrix algebra of invariants. Let  $X_r = (x_{ij}(r))_{1 \leq i, j \leq n}$  be the generic matrices of order  $n$  ( $r = \overline{1, d}$ ), and let  $\sigma_k(A)$  be the coefficients of the characteristic polynomial of a matrix  $A \in M_n(K)$ , that is  $\det(\lambda E - A) = \lambda^n - \sigma_1(A)\lambda^{n-1} + \cdots + (-1)^n \sigma_n(A)$ . The algebra  $R_{n,d}$  is generated by all elements of the form  $\sigma_k(X_{i_1} \cdots X_{i_s})$  (see [5]). The Procesi–Razmyslov Theorem on the relations in  $R_{n,d}$  was extended to the case of a field of an arbitrary characteristic in [14].

The goal of the constructive theory of invariants is to find a minimal (i.e. irreducible) homogeneous system of generators (shortly m.h.s.g.) for the algebra of invariants. A m.h.s.g. for  $R_{2,d}$  was determined in [12] for  $p = 0$ , in [10] for  $p > 2$ , and in [4] for  $p = 2$ . In [3] some upper and lower bounds on the highest degree of elements of a m.h.s.g. for  $R_{n,d}$  are pointed out for an arbitrary  $p$ . In [1] in the case  $p = 0$  the cardinality of a m.h.s.g. for  $R_{3,d}$  was calculated for  $d \leq 10$  on a computer, and was shown a way how such set can be constructed by means of a computer programme. The explicit upper bound on the highest degree of elements of a m.h.s.g. for  $R_{3,d}$  is given in [9] (except for the case  $p = 3, d = 6k + 1, k > 0$ , where the least upper bound is estimated with error not greater than 1). In this section we point out a m.h.s.g. for  $R_{3,d}$  for an arbitrary  $p, d$ .

The algebra  $R_{n,d}$  possesses natural  $\mathcal{N}$ - and  $\mathcal{N}^d$ -gradings by degrees and multide-

degrees respectively. Denote by  $R_{n,d}^+$  the subalgebra generated by all elements of  $R_{n,d}$  of positive degree. An element  $r \in R_{n,d}$  is called *decomposable*, if it can be expressed in terms of elements of  $R_{n,d}$  of lower degree, that is it belongs to the ideal  $(R_{n,d}^+)^2$ . Clearly,  $\{r_i\} \in R_{n,d}$  is a m.h.s.g. if and only if  $\{\bar{r}_i\}$  is a basis for  $\bar{R}_{n,d} = R_{n,d}/(R_{n,d}^+)^2$ . If two elements  $r_1, r_2 \in R_{n,d}$  are equal modulo the ideal  $(R_{n,d}^+)^2$ , we write  $r_1 \equiv r_2$ . There is a close connection between decomposability of an element of  $R_{n,d}$  and equality to zero of some element of  $N_{n,d}$  (see Lemma 18 below). Let  $A_{n,d}$  be a  $K$ -algebra without unity, generated by the generic matrices  $X_1, \dots, X_d$ . The homomorphism of algebras  $\Phi : A_{n,d} \rightarrow N_{n,d}$ , defined by  $\Phi(X_i) = x_i$ , is defined correctly (see [9]).

**Remark 4** For each  $\Delta = (\delta_1, \dots, \delta_d)$ , where  $\delta_1 \geq \dots \geq \delta_d \geq 0$ , let  $G_\Delta \subset R_{n,d}$  be such a set that its image in  $\bar{R}_{n,d}$  is a basis for  $\bar{R}_{n,d}(\Delta)$ . For any multidegree  $\Delta = (\delta_1, \dots, \delta_d)$  define  $G_\Delta$  by the following way:

$$G_\Delta = G_{\delta_{\sigma(1)}, \dots, \delta_{\sigma(d)}}|_{x_{ij}(r) \rightarrow x_{ij}(\sigma(r)), i, j = \overline{1, n}, r = \overline{1, d}},$$

where  $\sigma \in S_d$ ,  $\delta_{\sigma(1)} \geq \dots \geq \delta_{\sigma(d)}$ . Then, the set  $G = \cup_{\delta_1, \dots, \delta_d \geq 0} G_\Delta$  is a m.h.s.g. for  $R_{n,d}$ .

Further, we assume that  $n = 3$ , unless it is stated otherwise.

Let  $B_\Delta$  be the basis for  $N_{3,d}(\Delta)$  from Proposition 2 and Theorems 2, 3. For  $u \in K\langle F_d \rangle^\#$  denote  $\text{tr}(u) = \text{tr}(u|_{x_i \rightarrow X_i, i = \overline{1, d}}) \in R_{3,d}$ .

**Theorem 4** For multidegree  $\Delta = (\delta_1, \dots, \delta_d)$ , where  $\delta_1 \geq \dots \geq \delta_d$ ,  $d \geq 1$ , define  $G_\Delta \subset R_{3,d}$ :

1) the case of  $p \neq 3$ :

if  $d \geq 2$  and  $\delta_d = 1$ , then  $G_\Delta = \{\text{tr}(ux_d) | u \in B_{(\delta_1, \dots, \delta_{d-1})}\}$ ,

if  $\Delta = 2^3$ , then  $G_\Delta = \{\text{tr}(X_1^2 X_2^2 X_3^2)\}$ ,

if  $\Delta = 2^2$ , then  $G_\Delta = \{\text{tr}(X_1^2 X_2^2)\}$ ,

if  $\Delta = 3^2$ , then  $G_\Delta = \{\text{tr}(X_1^2 X_2^2 X_1 X_2)\}$ ,

if  $d = 1$ ,  $\Delta = k$ ,  $k = \overline{1, 3}$ , then  $G_\Delta = \{\sigma_k(X_1)\}$ ,

for others  $\Delta$  we define  $G_\Delta = \emptyset$ ;

2) the case of  $p = 3$ :

if  $d \geq 2$  and  $\delta_d = 1, 2$ , then  $G_\Delta = \{\text{tr}(ux_d^{\delta_d}) | u \in B_{(\delta_1, \dots, \delta_{d-1})}\}$ ,

if  $\Delta = 3^{2k}$ ,  $k > 0$ , or  $\Delta = 3^{6k+1}$ ,  $k > 0$ , then  $G_\Delta = \{\text{tr}(u) | u \in B_\Delta\}$ ,

if  $d = 1$ ,  $\Delta = k$ ,  $k = \overline{1, 3}$ , then  $G_\Delta = \{\sigma_k(X_1)\}$ ,

for others  $\Delta$  we define  $G_\Delta = \emptyset$ .

Then, the set  $G$  from Remark 4 is a minimal system of generators for  $R_{3,d}$ .



In order to prove the Theorem we need some statements from [9] and its corollaries:

**Lemma 18** 1. Let  $H \in A_{3,d-1}$ . Then  $\text{tr}(HX_d)$  is decomposable if and only if  $\Phi(H) = 0$  in  $N_{3,d}$ .

2. If  $\text{tr}(HX_d^2) \equiv 0$ , where  $H \in A_{3,d-1}$ , then  $\Phi(H)x_d + x_d\Phi(H) = 0$  in  $N_{3,d}$ .

3. Let  $p = 3$ ,  $H \in A_{3,d-1}$ . Then  $\text{tr}(HX_d^2) \equiv 0$  if and only if  $\Phi(H) = 0$  in  $N_{3,d-1}$ .

4. If  $\text{tr}(u)$  is indecomposable, where  $u$  is a word, then there are canonical words  $u_i$ ,  $\text{mdeg}(u) = \text{mdeg}(u_i)$ , and  $\alpha_i \in K$ , such that  $\text{tr}(u) \equiv \sum \alpha_i \text{tr}(u_i)$ .

5. We have  $\sigma_2(UV) \equiv \text{tr}(U^2V^2)$ , where  $U, V \in A_{3,d}$ .

6. Elements  $\sigma_2(X_1)$ ,  $\det(X_1)$  are indecomposable.

7. If  $p \neq 3$ , then  $\text{tr}(X_1^2X_2^2X_3^2) + \text{tr}(X_1^2X_3^2X_2^2) \equiv 0$ . For any  $p$ , the element  $\text{tr}(X_1^2X_2^2X_3^2)$  is indecomposable.

8. Let  $p = 3$ ,  $u_i \in F_d$  are words,  $\text{mdeg}(u_i) = \Delta$ . Then  $\sum_i \alpha_i \text{tr}(u_i) \equiv 0$  if and only if  $\sum_i \alpha_i u_i = 0$  is a consequence of the system of identities  $\mathcal{S}_\Delta$  and identities  $uv = vu$ , where  $u, v \in F_d^\#$  and  $\text{mdeg}(uv) = \Delta$ .

9. For  $u, v \in F_d$ , where  $\text{mdeg}(uv) = 3^{2k}$ ,  $k > 0$ , or  $\text{mdeg}(uv) = 3^{6k+1}$ ,  $k > 0$ , we have  $uv = vu$  in  $N_{3,d}$ .

10. Let  $D$  be the explicit upper bound on degrees of elements of a m.h.s.g. for  $R_{3,d}$ , where  $d \geq 2$ . Then

if  $p = 0$  or  $p > 3$ , then  $D = 6$ ;

if  $p = 2$ , then  $D = \begin{cases} d+2 & , \quad d \geq 4 \\ 6 & , \quad d = 2, 3 \end{cases}$ ;

if  $p = 3$  and  $d = 6k + r$ , where  $r \in \{3, 5\}$ ,  $k \geq 0$ , then  $D = 3d - 1$ .

**Proof.** All items except for 3, 7 are proven in [9].

3. If  $\text{tr}(HX_d^2) \equiv 0$ , then  $\Phi(H)x_d + x_d\Phi(H) = 0$  in  $N_{3,d}$  (see item 2). By item 4 of Lemma 5 we have  $2\Phi(H) = 0$  in  $N_{3,d}$ . The converse follows from item 1.

7. Let  $p \neq 3$ . The identity  $x_1^2x_2x_3^2 = 0$  in  $N_{3,d}$  (Lemma 1 item 4) implies  $\text{tr}(X_1^2X_2X_3^2X_2) \equiv 0$  (see item 1). On the other hand, the identity  $x_2x_3^2x_2 = -x_2^2x_3^2 - x_3^2x_2^2$  in  $N_{3,d}$  (see identity (1)) implies  $\text{tr}(X_1^2X_2X_3^2X_2) \equiv -\text{tr}(X_1^2X_2^2X_3^2) - \text{tr}(X_1^2X_3^2X_2^2)$  (see item 1). The claim is proved.

Assuming  $\text{tr}(X_1^2X_2^2X_3^2) \equiv 0$ , we get  $x_1^2x_2^2x_3 + x_3x_1^2x_2^2 = 0$  in  $N_{3,d}$  by item 2. Thus  $x_1^2x_2^2x_1 = -x_1x_1^2x_2^2 = 0$  in  $N_{3,d}$ ; that is a contradiction to item 2 of Lemma 1.  $\triangle$

**Proof of theorem 4.** By items 4, 5 of Lemma 18, we have that  $R_{3,d}$  is generated by  $\{\sigma_2(X_i), \det(X_i), \text{tr}(u) \mid u \in F_d^\# \text{ is a canonical word, } i = \overline{1, d}\}$ .

Let  $p \neq 3$ . The claim follows from items 1, 6, 7 and 10 of Lemma 18.

Let  $p = 3$ . The case of  $\delta_d = 1, 2$  follows from items 1, 3 of Lemma 18. The case of  $\Delta = 3^d$  follows from items 8, 9, 10 of Lemma 18.  $\triangle$

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